

# Distributed Systems

## 16. Naming

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# Naming things

- Naming: map names to objects
  - Helps with using, sharing, and communicating information
- Examples
  - **User names**: *used for system login, email, chat*
  - **Machine names**: *used for ssh, email, web*
  - Files
  - Devices
  - Objects, functions, variables in programs
  - Network services

# What's a name?

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**Name:** identifies what you want

**Address:** identifies where it is

**Route:** identifies how to get there

**Binding:** the association of a name with the object

- “choose a lower-level-implementation for a higher-level semantic construct”

*RFC 1498: Inter-network Naming, addresses, routing*

# Pure & Impure Names

- **Pure names** – *identify*
  - The name contains no information aside from the name
  - It does not identify *where* the object can be found
  - Examples:
    - `c8:2a:14:3f:92:d1` my computer's ethernet MAC address
    - `paul.krzyzanowski` my facebook name
    - `908-555-3836` phone # (this used to be an impure name)

# Pure & Impure Names

- **Impure names** – *guide*
  - The name contains context information
  - Object is generally unmovable
  - Examples:
    - `p@pk.org`, `pxk@cs.rutgers.edu`
      - User names in different Internet domains: same person or not?
      - Context (domain name) is encoded into the name
    - `/home/paul/bin/qsync`
      - File pathname changes if we move the object

# Uniqueness of names

- Easy on a small scale – problematic on a large scale
  - It can be difficult to make globally unique names
- **Uniqueness for pure names**
  - Designate a bit pattern or naming prefix that does not convey information
    - **Ethernet MAC address**: 3 bytes: organization, 3 bytes: controller
    - **IP address**: network & host (variable partition)
- **Uniqueness for impure names**
  - Use a **hierarchy**
  - Globally unique components (pure names)
    - **Compound name**: iterative list of pure names connected with separators
      - Domain name ([www.cs.rutgers.edu](http://www.cs.rutgers.edu))
      - URLs (<http://pk.org/417/lectures/l-naming.html>)
      - File pathnames (</usr/share/dict/words>)

# Terms: Naming convention = syntax

Naming system determines syntax for names

Naming convention can take any format

- Ideally one that will suit the application and user
  - E.g., human readable names for humans, binary identifiers for machines
- UNIX file names:
  - Parse components from left to right separated by /  
`/home/paul/src/gps/gui.c`
- Internet domain names:
  - Ordered right to left and delimited by .  
`www.cs.rutgers.edu`
- LDAP names
  - Attribute/value pairs ordered right to left, delimited by ,  
`cn=Paul Krzyzanowski, o=Rutgers, c=US`

# Terms: Context

A particular set of *name* → *object* bindings

- Names are unique within the context
  - E.g., `/etc/postfix/main.cf` on a specific computer
- Each context has an associated naming convention
- A name is always interpreted relative to some context
  - E.g., directory `/usr` in a Linux file system on `crapper.pk.org`



# Terms: Naming System

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Connected set of contexts of the same type (same naming convention) along with a common set of operations

For example:

- System that implements DNS (Internet domain names)
- System that implements LDAP (directory of people)

# Terms: Namespace = set of names

- A container for a set of names in the naming system
- A namespace has a scope
  - **scope** = region where the name exists & refers to the object
  - For example,
    - Names of all files in a directory
    - All domain names within rutgers.edu
    - E.g., Java package, local variables
- A namespace may be tree structured (hierarchical)
  - Fully-qualified or hierarchical names may be used to identify names outside the local namespace
  - Global namespace = root of the tree

# Terms: Resolution

- Resolution = name lookup
  - Return the underlying representation of the name
  - Look up the **binding** of the name to its object
- For example,
  - [www.rutgers.edu](http://www.rutgers.edu) → 128.6.4.5
- **Iterative** resolution
  - Example: parse a pathname
- **Recursive** resolution
  - Example: parse a distribution list: each entity may be expanded

# When do should you do a resolution?

## Static binding

- Hard-coded

## Early binding

- Look up binding before use
- Cache previously used binding

## Late binding

- Look up just before use



These can cause problems!

# Name Service

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The service that performs name resolution

Allows you to resolve **names**

- Looking up a **name** gives the corresponding **address** as a response

Can be implemented as

- Search through file
- Database query
- Client-server program (**name server**) – may be distributed
- ...

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# Directory Service

- Extension of name service:
  - Associates names with objects
  - Allows objects to have attributes
  - Can search based on attributes
  
- For example,
  - LDAP (Lightweight Directory Access Protocol)
  - Directory can be an object store:
    - E.g., look up printer object and send data stream to it

# IP Domain Names

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Human readable names

e.g., [www.cs.rutgers.edu](http://www.cs.rutgers.edu)

Hierarchical naming scheme

- Top of hierarchy on the right
- No relation to IP address or network class

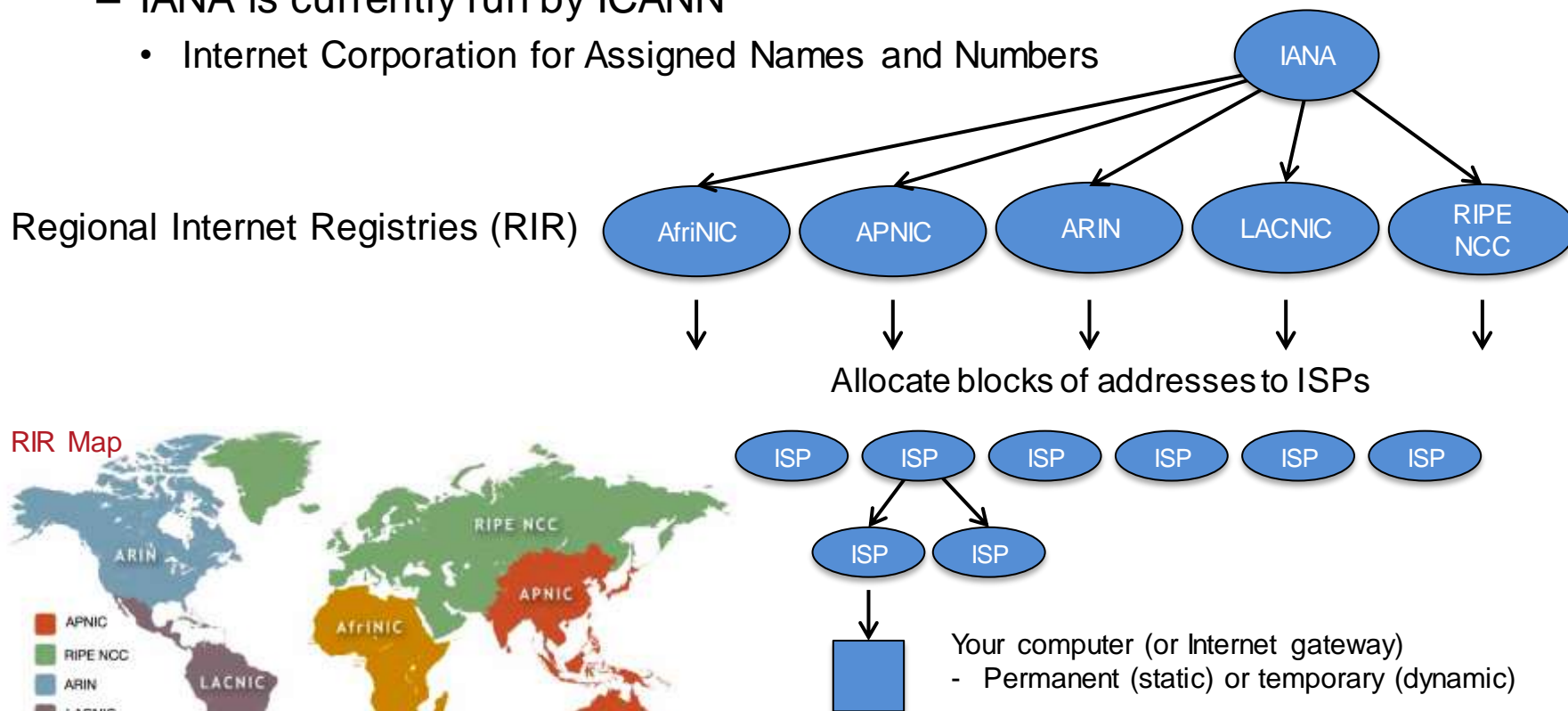


# Case Study: Internet Domain Name System (DNS)

# How are IP addresses assigned?

IP addresses are distributed hierarchically

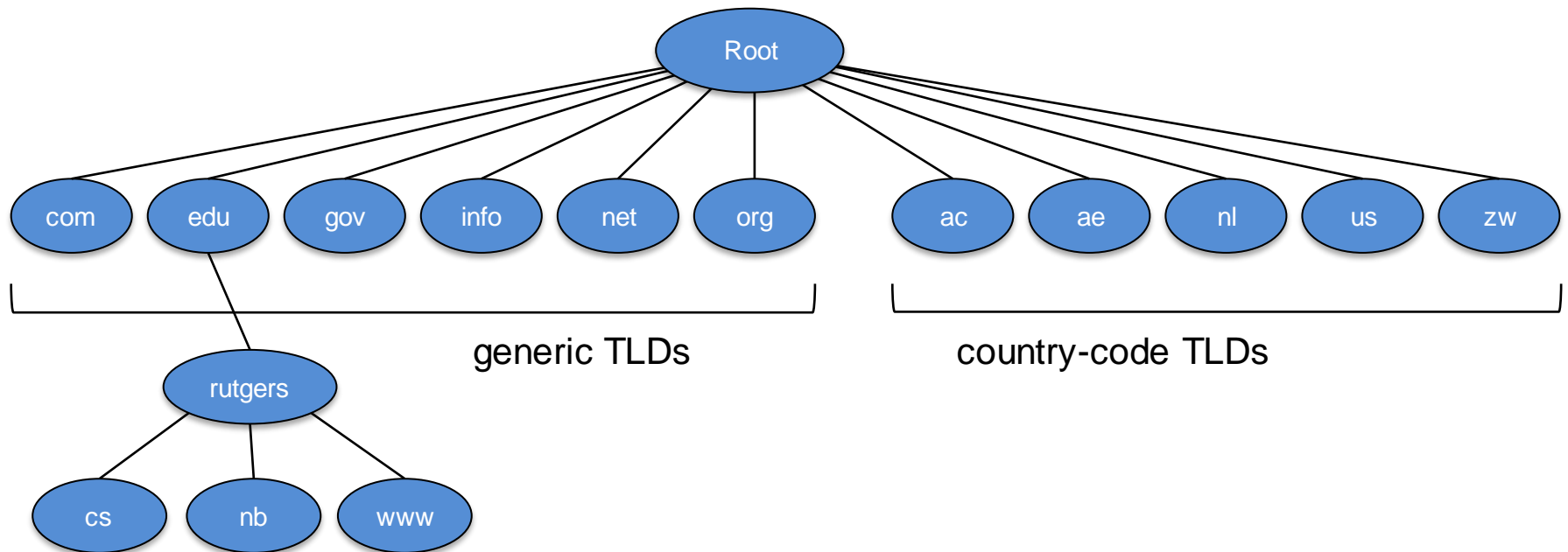
- Internet Assigned Numbers Authority (IANA) at the top
  - IANA is currently run by ICANN
    - Internet Corporation for Assigned Names and Numbers



# How are machine names assigned?

- Early ARPANET
  - Globally unique names per machine (e.g., UCBVAX)
  - Kept track at the Network Information Center (NIC) at the Stanford Research Institute (SRI)
- That doesn't scale!
- A **domain hierarchy** was created in 1984 (RFC 920)
  - Domains are administrative entities: divide name management
  - Tree-structured global name space
  - Textual representation of domain names  
`www.cs.rutgers.edu`

# Domain Name Hierarchy



# Top Level Domains (TLDs)

## ccTLD

**Country-code** domains  
ISO 3166 codes  
e.g., .us, .de, .ca, .es

## IDN ccTLD

**Internationalized**  
country-code domains  
e.g., .السعودية, .中國, .php

## gTLD

**Generic** top-level domains  
e.g., .biz, .com, .edu,  
.gov, .info, .net, .org,  
.audio, .catering, .网络

There are currently 1,446 top-level domains (as of Oct 31, 2016)

Each top-level domain has an administrator assigned to it

Assignment is delegated to various organizations by the Internet Assigned Numbers Authority (IANA)

IANA keeps track of the **root servers**

See <http://www.iana.org/domains/root/db> for the latest count

# Shared registration

- **Domain name registry:** *this is the database*
  - Keeps track of all domain names registered under a top-level domain
- **Domain name registry operator:** *this is the company that runs the DB*
  - NIC = **Network Information Center** – organization that keeps track of the registration of domain names under a top-level domain
    - Keeps the database of domain names
    - See <https://www.icann.org/resources/pages/listing-2012-02-25-en>
- **Domain name registrar:** *this is the company you use to register*
  - Company that lets you register a domain name
  - Registrars update the registry database at the NIC

# Shared registration

- Multiple domain **registrars** provide domain **registration services**
  - 2,147 registrars as of October 2016, including 701 unique DropCatch.com registrars
- The registrar you choose becomes the **designated registrar** for your domain
  - Maximum period of registration for a domain name = 10 years
- The **registry operator** keeps the **central registry database** for the top-level domain
- Only the designated registrar can change information about domain names
  - A domain name owner may invoke a domain transfer process

## Example

- *Namecheap* is the designated registrar for *poopybrain.com*
- *VeriSign, Inc.* is the registry operator for the *.com* gTLD

See <https://www.icann.org/registrar-reports/accredited-list.html> for the latest list of registrars

# The problem

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Every device connected to the internet has a unique Internet Protocol (IP) address

How do you **resolve** user-friendly machine names to IP addresses?

www.cs.rutgers.edu → 128.6.4.24



# Original solution

## Through the 1980s

- Search `/etc/hosts` file for machine name (see RFC 606)
- File periodically downloaded from Network Information Center (NIC) at the Stanford Research Institute (SRI)
- This was not sustainable with millions of hosts on the Internet
  - A lot of data
  - A lot of churn in the data
    - new hosts added, deleted, addresses changed
  - Maintenance
  - Traffic volume

*Solution doesn't scale!*

# DNS: Domain Name System

- Distributed database: a hierarchy of **name servers**
- **DNS** is an application-layer protocol
  - Name-address resolution is handled at the edge
  - The network core is unaware of host names ... and does not care
  - There is no special relationship between names and addresses
    - Example: **cs.poopybrain.com** can resolve to **cs.rutgers.edu**

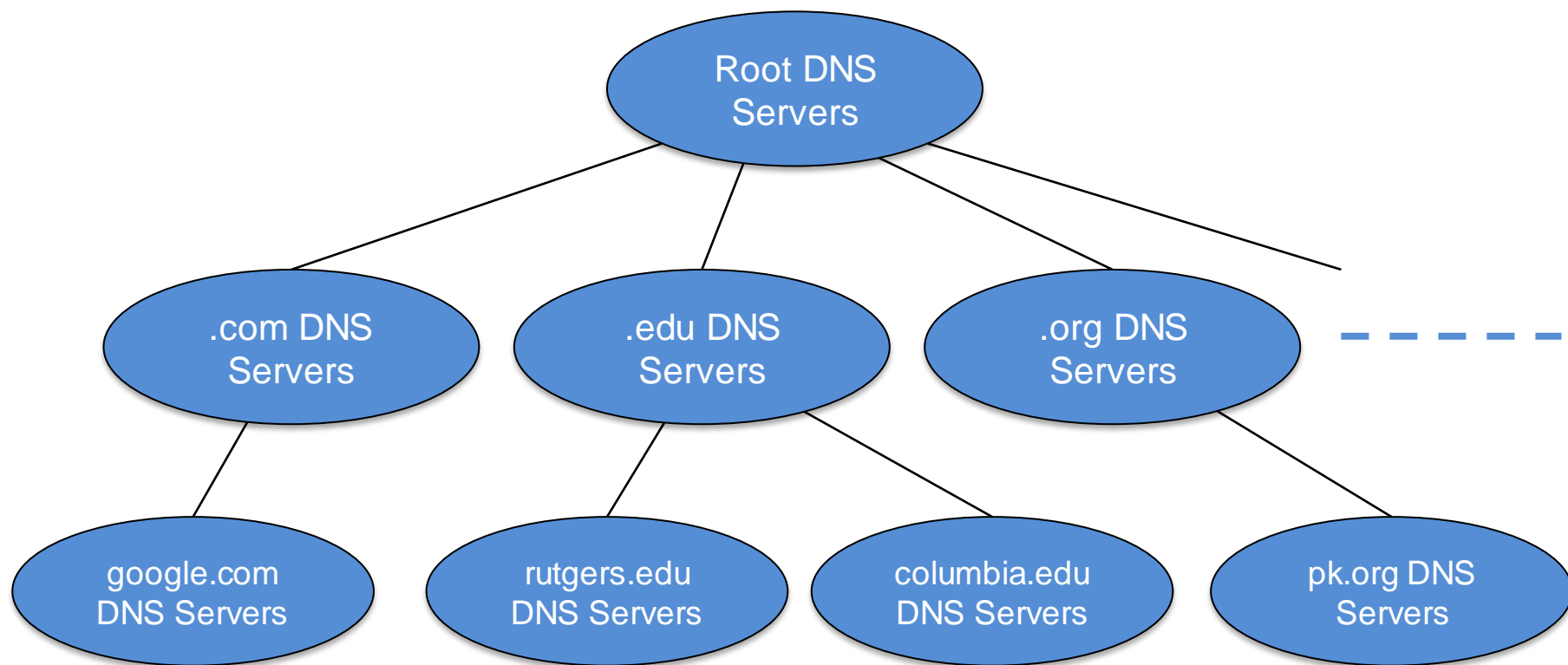
cs.poopybrain.com → cs.rutgers.edu

# DNS provides

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- Name to IP address translation
- Aliasing of names (called **canonical** names)
- Identification of name servers
- Mail server names
- Load distribution:
  - Multiple name servers may handle a query for a domain
  - Caching – store past look-ups
  - Ability to provide a set of IP addresses for a name

# DNS is a distributed, hierarchical database



A collection of DNS servers

# Authoritative DNS server

- An **authoritative name server** is responsible for answering queries about its zone
  - Provides *real* answers vs. *cached* answers
  - Configured by the administrator
- **Zone** = group of machines under a node in the tree  
E.g., rutgers.edu

# A DNS server returns answers to queries

Key data that a DNS server maintains (partial list)

Information	Abbreviation	Description
Host	A	Host address (name to address) Includes name, IP address, time-to-live (TTL)
Canonical name	CNAME	Name for an alias
Mail exchanger	MX	Host that handles email for the domain
Name server	NS	Identifies the name server for the zone: tell other servers that yours is the authority for info within the domain
Start of Zone Authority	SOA	Specifies authoritative server for the zone. Identifies the zone, time-to-live, and primary name server for the zone

# Finding your way

- How do you find the DNS Server for rutgers.edu?
  - That's what the **domain registry** keeps track of
  - When you register a domain,
    - You supply the addresses of at least two DNS servers that can answer queries for your zone
    - You give this to the **domain registrar**, who updates the database at the **domain registry**
- So how do you find the right DNS server?
  - Start at the root

# Root name servers

- The **root name server** answers can return a list of authoritative name servers for top-level domains
- 13 root name servers
  - A.ROOT-SERVERS.NET, B.ROOT-SERVERS.NET, ...
  - Each has redundancy (via *anycast* routing or load balancing)
    - Each server is really a set of machines



Download the latest list at <http://www.internic.net/domain/named.root>



# DNS Queries

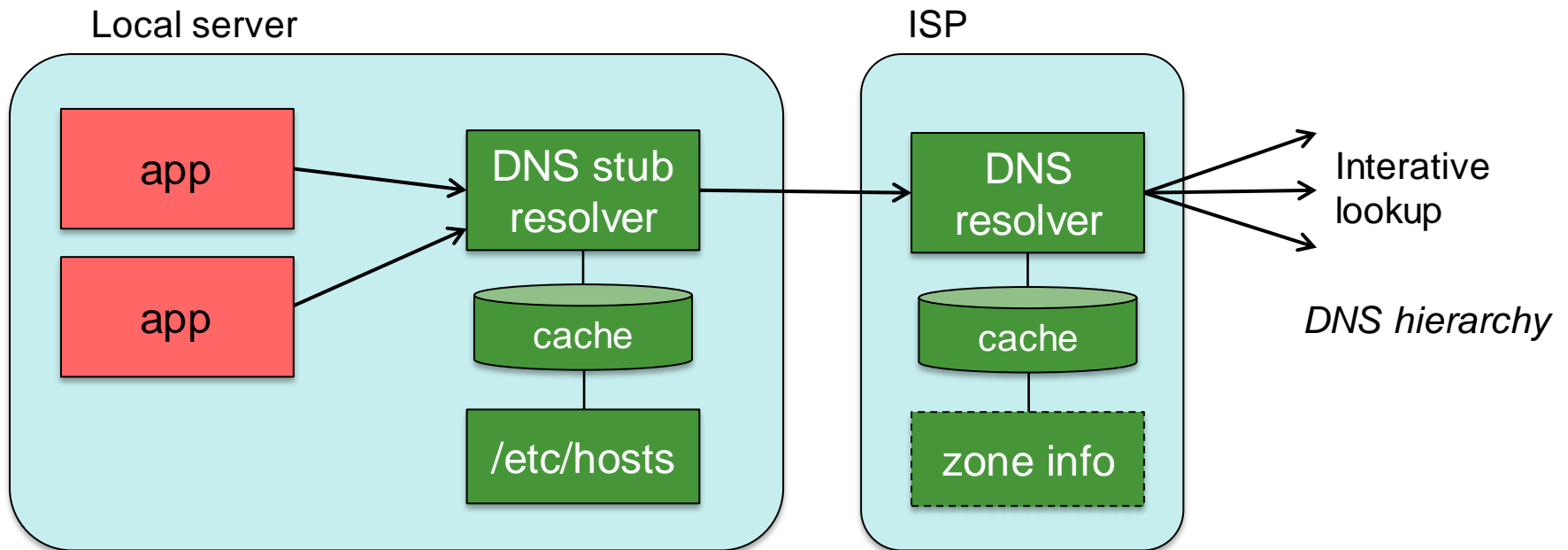
- **Iterative** (non-recursive) name resolution
  - DNS server will return a definitive answer or a **referral** to another DNS server
    - **referral** = reference to a DNS server for a lower level of the queried namespace
    - Server returns intermediate results to the client
      1. Send query to a **root** name server
      2. Send query to a **edu** name server
      3. Send query to a **rutgers** name server
  - **Advantage: stateless**
- **Recursive** DNS name resolution
  - Name server will take on the responsibility of fully resolving the name
    - May query multiple other DNS servers on your behalf
  - **DNS server cannot refer the client to a different server**
  - **Disadvantage: name server has more work; has to keep track of state**
  - **Advantages: Caching opportunities, less work for the client!**

*Most top-level DNS servers only support iterative queries*

# DNS Resolvers: local name server

- **DNS Resolver** = client side of DNS
  - Not really a part of the DNS hierarchy
  - Acts as an intermediary between programs that need to resolve names and the name servers
  - A resolver is responsible for performing the full resolution of the query
- Where are the resolvers?
  - Each local system has one: that's what applications contact
    - Local cache; may be a process or a library
    - On Linux & Windows, these are limited DNS servers (called **stub resolvers**)
      - Usually not capable of handling referrals and expect to talk with a name server that can handle recursion (full resolution)
  - ISPs (and organizations) run them on behalf of their customers
    - Including a bunch of free ones (OpenDNS, Google Public DNS)
- Resolvers cache past lookups – they are not responsible for zones

# DNS Resolvers in action



## Local stub resolver:

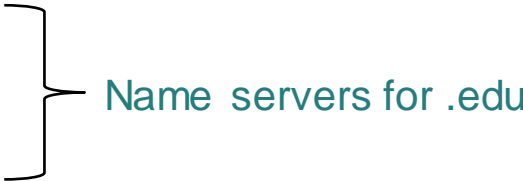
- check local cache
- check local hosts file
- send request to external resolver

E.g., on Linux: resolver is configured via the `/etc/resolv.conf` file

## External resolver

- DNS server that accepts recursion
- Running at ISP, Google Public DNS, OpenDNS, etc.

# Sample query

- Rutgers registered rutgers.edu with the .edu domain
  - educause.net is the domain registry for the .edu gTLD
  - Registration includes defining the name servers for .rutgers.edu
    - ns124.a2.incapsecuredns.net: 192.230.123.124
    - ns8.a1.incapsecuredns.net: 192.230.122.8
    - ns87.a0.incapsecuredns.net: 192.230.121.87
- EDUCAUSE registered its name servers with root name servers
  - ns1.twtelecom.net
  - ns1.educause.edu
  - ns1.twtelecom.net

Name servers for .edu
- We know how to get to root name servers
  - Download <http://www.internic.net/domain/named.root>

# Sample Query

Submit query to a local *DNS resolver*:

1. *query(cs.rutgers.edu)* → any root name server  
send query to c.root-servers.net: 192.33.4.12
2. Receive *referral* to a list of DNS servers for *edu*  
a.edu-servers.net: 192.5.6.30                      g.edu-servers.net: 192.42.93.30
3. *query(cs.rutgers.edu)* → edu name server  
send query to g.edu-servers.net: 192.41.162.32
4. Receive *referral to rutgers.edu* name servers:
  - ns87.a0.incapsecuredns.net      192.230.121.86
  - ns8.a1.incapsecuredns.net.192.230.122.7
  - ns124.a2.incapsecuredns.net      192.230.123.123
5. *query(cs.rutgers.edu)* → rutgers name server  
send query to 192.230.122.7
6. The rutgers name server returns
  - A: 128.6.4.2                      *address*
  - MX: dragon.rutgers.edu            *domain name for email*

# Caching

- Starting every query at the root would place a huge load on root name servers
- A name server can cache results of previous queries
  - Save query results for a *time-to-live* amount of time
  - The time-to-live value is specified in the domain name record by an authoritative name server

# The End