Distributed Systems

Fall 2017 Exam 3 Review

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The core task of the *user's map function* within a *map* worker in a MapReduce framework is to:

- (a) Determine which reduce worker should process which key.
- (b) Split the input data into shards.
- (c) Parse input data and create key, value tuples.
- (d) All of the above.

Framework – splits data

Partitioning function – determines which reduce worker handles a key

In MapReduce, partitioning refers to:

- (a) Determining the ratio of map workers to reduce workers.
- (b) Determining which reduce worker will process a specific key.
- (c) Splitting the input data into shards.
- (d) Assigning input shards to map workers

Reduce workers in MapReduce can start working:

- (a) In parallel when the map workers start.
- (b) When at least one map worker starts to generate data.
- (c) When at least one map worker has processed all its input.
- (d) When every single map worker has completed its task

<u>All</u> <key, value> sets must be generated before <u>any</u> reducer can start

Bigtable's *multidimensional* property refers to the fact that:

- (a) Bigtable stores versioned data within rows and columns.
- (b) A table is actually composed of an arbitrary number of tablets.
- (c) A multi-level storage structure is used: memtable, SSTable, tablet, and table.
- (d) Each cell in a table can also be a table and, recursively, cells within that table can be tables.

d. Not supported in Bigtable

As new rows are added to a Bigtable, they are:

- (a) Added to an arbitrary tablet in the table that has free space.
- (b) Appended to the end of the entire table.
- (c) Appended to the end of the entire table but an index file with sorted keys enables rapid lookup.
- (d) Added in a way to make sure the table remains sorted by a single key.

Tablets & tables are always kept sorted.

In Bigtable, what is the unit of distribution and load balancing?

- (a) A set of adjacent rows.
- (b) A set of adjacent columns.
- (c) Each column family.
- (d) Timestamped versions of data.

Tablets are broken along rows.

To coordinate transaction commits across multiple servers, Spanner uses:

- (a) A two-phase commit protocol.
- (b) A three-phase commit protocol.
- (c) Distributed consensus based on Paxos.
- (d) Optimistic concurrency control, checking for problems after the commit.

To provide *isolation* of transactions, Spanner:

- (a) Restricts execution to one transaction at a time.
- (b) Uses two-phase locking.
- (c) Uses strict two-phase locking.
- (d) Requires transactions to specify the data they plan to access ahead of time.

TrueTime provides:
(a) A means of synchronizing clocks across multiple data centers.
(b) A bounded time interval that contains the actual time of day within the interval.
(c) The exact time of day obtained from local time servers.
(d) A vector clock to enable each transaction to obtain a unique time stamp.

- a. Each data center is responsible for its own clock synchronization and has its own master clocks: GPS & an atomic clock
- c. Synchronization algorithms never give us the exact time. TrueTime supplies a range.

Spanner addresses the problem of global time ordering by:

- (a) Allowing each transaction to get the precise time of day.
- (b) Using consistent (total) ordering instead of global time ordering.
- (c) Using an eventual consistency model where time of day does not matter.
- (d) Forcing commit operations to wait.

Commit wait = wait until the timestamp of the transaction is definitely in the past.

Spanner allows transactions to use lock-free reads by:

- (a) Using optimistic concurrency control mechanisms and not using write locks.
- (b) Letting them read from replicas instead of the main servers.
- (c) Using write locks but no read locks
- (d) Letting them read older versions of data.

Spanner implements multiversion concurrency.

Old versions of data are readable while transactions are modifying new data. Other transactions can see a consistent, but slightly older, view of the world.

Messages sent by a process during execution of a superstep in BSP:

- (a) Must be delivered before the start of the next superstep.
- (b) Are delivered only at the start of the next superstep.
- (c) Can be delivered to any programmer-specified future superstep.
- (d) Are multicast to the entire group and acknowledged at the end of the superstep.

End of superstep = barrier

In Pregel, a function is executed for:

- (a) Each vertex of a graph.
- (b) Each edge of a graph.
- (c) A graph cluster, representing a connected set of vertices and their edges.
- (d) Each subgraph that is allocated to a distinct server.

Pregel's combiners:

- (a) Reduce the number messages from the same processor that are targeted to the same destination.
- (b) Manage global state.
- (c) Merge multiple vertices into one vertex.
- (d) Merge multiple edges into one edge.

Combiner = optional function to consolidate messages to the same vertex

Aggregator = global state

- In Spark, a *Resilient Distributed Dataset*, or RDD, is:
- (a) A distributed collection of objects that is modified by each transformation.
- (b) An immutable distributed collection of objects representing original data or the output of a transformation.
- (c) The original input data that will be processed by Spark and is replicated onto multiple servers.
- (d) The output data generated by a Spark action.
- a. RDD immutable = never modified
- c. RDD can be original data or the output of a transformation
- d. Actions produce final data

Spark's *fault tolerance* is based on:

- (a) Checkpointing the output of each transformation and action.
- (b) Running replicated transformation servers.
- (c) Keeping track of the sequence of transformations that created the needed data.
- (d) Restarting the entire sequence of transformations from the user's original data.

Spark backtracks to try get the latest available RDDs.

Multihoming means:

- (a) A process migrates between multiple servers.
- (b) Content is cached in multiple places close to the user.
- (c) A system is connected to more than one network.
- (d) The same content may be generated from multiple sources.

Akamai's *dynamic DNS* (domain name service):
(a) Locates the most suitable edge server based on a client's URL request.
(b) Locates the most suitable edge server based on a client's domain name query.
(c) Locates the shortest path to the origin server from a specific client.
(d) Locates the set of edge servers that should cache content for a specific host.

a. DNS doesn't see URL requests

- c. The transport network handles the shortest path
- d. Dynamic DNS doesn't give a list of servers for caching content.

A system area network is typically designed to:

- (a) Eliminate the overhead of TCP while providing reliable communication.
- (b) Be a dedicated network for storage components.
- (c) Act as a heartbeat network to allow detection of network failures.
- (d) Connect hardware elements within a computer system.

A *clustered file system* differs from a distributed file system in that:

- (a) Multiple computers access the same physical storage device.
- (b) Data may be distributed among multiple computers.
- (c) Data is replicated across storage devices on multiple computers for fault tolerance.
- (d) It provides services only over a local area network.

A clustered file system does NOT:

- (a) Require a distributed lock manager.
- (b) Access data on a device block level rather than a file level.
- (c) Enable multiple systems to share files.
- (d) Distribute a file's data among multiple servers.

a. Because storage devices are shared, a distributed lock manager is required.

b. By definition, clustered file systems read & write raw blocks.

c. Clustered file systems are designed to provide concurrent access from multiple systems.

Fencing is used to:

- (a) Provide a trusted path for nodes to communicate on a LAN.
- (b) Isolate a computing node from other nodes.
- (c) Monitor whether cluster members are alive.
- (d) Establish a quorum among cluster members.

Fencing shuts off or isolates components that may be misbehaving.

Unlike a public key algorithm, a symmetric algorithm:

- (a) Uses the same function for encryption as decryption.
- (b) Uses the same key for encryption and decryption.
- (c) Produces ciphertext that is the same length as the plaintext.
- (d) Cannot be used for message authentication.

For Alice to send an *encrypted signed* message to Bob, she creates a hash of the message and sends Bob:

- (a) The message encrypted with Alice's private key and the hash encrypted with Bob's public key.
- (b) The message encrypted with Alice's public key and the hash encrypted with Alice's private key.
- (c) The message encrypted with Bob's public key and the hash encrypted with Alice's private key.
- (d) The message encrypted with Bob's public key and the hash encrypted with Alice's public key.

A message encrypted with Bob's public key can only be decrypted by Bob.

A hash encrypted with Alice's private key could have been encrypted only be Alice.

A cryptographic hash function is an example of a:

- (a) One-way function.
- (b) Message authentication code.
- (c) Symmetric algorithm.
- (d) Session key.

(b) A MAC is an encrypted hash of a message.

(d) This is just a random number.

The *Diffie-Hellman* algorithm most directly solves the problem of:

- (a) Alice being able to send authenticated messages to Bob.
- (b) Alice being able to validate Bob's identity.
- (c) Alice and Bob generating public keys.
- (d) Alice and Bob getting a shared secret key.

The Diffie-Hellman algorithm was created for key exchange.

The Diffie-Hellman algorithm is not needed if you have:

- (a) Hash functions.
- (b) Message authentication codes.
- (c) Symmetric cryptography.
- (d) Public key cryptography

(a) & (b) – do not facilitate key exchange

(c) On its own, does not enable key exchange: need a trusted 3rd party

Salt in a password hash is used to:

- (a) Implement single-use (one-time) passwords.
- (b) Add a layer of protection against bad hash functions.
- (c) Encrypt the password before generating the hash.
- (d) Make attacks using precomputed hash tables ineffective.

Salt is randomly-generated – but not secret – junk appended to the password before it is hashed.

Linux /etc/shadow entry:

poopybrain:\$6\$7oRkRWSd\$d9GJSs8tMUdg6LrbwzeocjKpCHpA/3dR/knwV/jkA/l8ZZIJNU63Tw3c35XJAIVqz5C5EEE4STn59mu.quiJv1:17511:0:99999:7:::

\$6\$ = SHA512 hash \$70RkRWSd\$ = Salt

\$d9G...JV1\$ = sha1_hash("monkey", salt)

- 29. An advantage of the Challenge-Handshake Authentication Protocol (CHAP) is:
- (a) The user or client does not need to know any secret information.
- (b) It is a time-based protocol and the password is invalid after a short time.
- (c) It does not require the use of one-way functions.
- (d) No secret information is sent on the network.

(a) Both sides need to know a secret.

(b) No.

(c) The response is hash(secret, challenge)

Kerberos is designed to allow Alice and Bob to communicate using:

- (a) A public key algorithm.
- (b) A symmetric cryptography algorithm.
- (c) A hybrid cryptosystem.
- (d) A restricted cipher.

Kerberos uses only symmetric cryptography.

Secure Sockets Layer (SSL, or Transport Layer Security, TLS) uses:

- (a) A public key algorithm.
- (b) A symmetric cryptography algorithm.
- (c) A hybrid cryptosystem.
- (d) A restricted cipher

SSL uses public key cryptography for key exchange (and authentication) and symmetric cryptography for communication.

OAuth was designed to:

- (a) Allow a user to grant one service specific access rights from another service.
- (b) Authenticate users using X.509 digital certificates.
- (c) Enable an administrator to authorize user access to services.
- (d) Support multi-factor authentication protocols.

Authentication mechanisms are not specified in OAuth. It's up to the service.

OAuth relies on:

- (a) HTTP URL redirection.
- (b) Public key cryptography.
- (c) A trusted third party that stores all the keys.
- (d) Kerberos to authenticate and authorize users.

The end