

CROW: A Low-Cost Substrate for Improving DRAM Performance, Energy Efficiency, and Reliability

Hasan Hassan Minesh Patel Jeremie S. Kim A. Giray Yaglikci

Nandita Vijaykumar Nika Mansouri Ghiasi Saugata Ghose Onur Mutlu



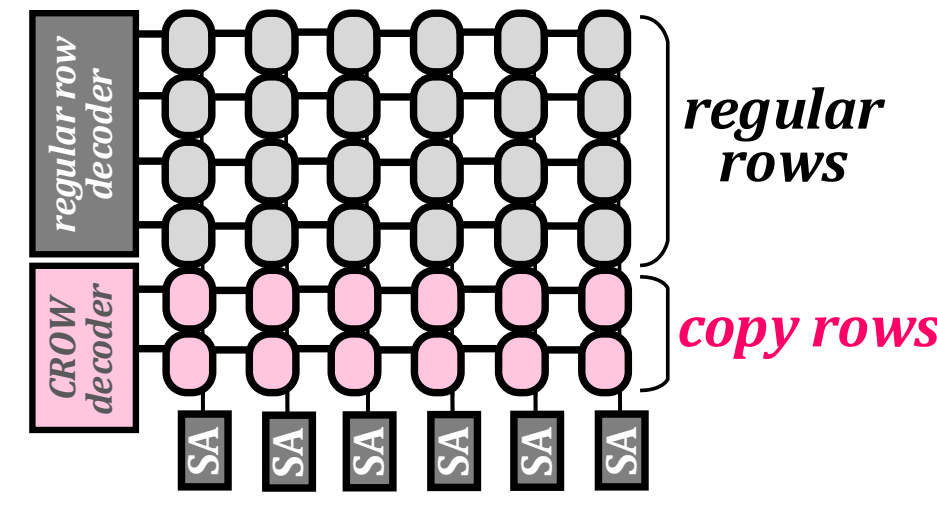
1: Summary

Challenges of DRAM scaling:

- **High access latency** → bottleneck for improving system performance/energy
- **Refresh overhead** → reduces performance and consume high energy
- **Exposure to vulnerabilities** (e.g., RowHammer)

Copy-Row DRAM (CROW)

- Introduces **copy rows** into a subarray
- The benefits of a **copy row**:
 - Efficiently duplicating data from regular row to a **copy row**
 - Quick access to a duplicated row
 - Remapping a regular row to a **copy row**

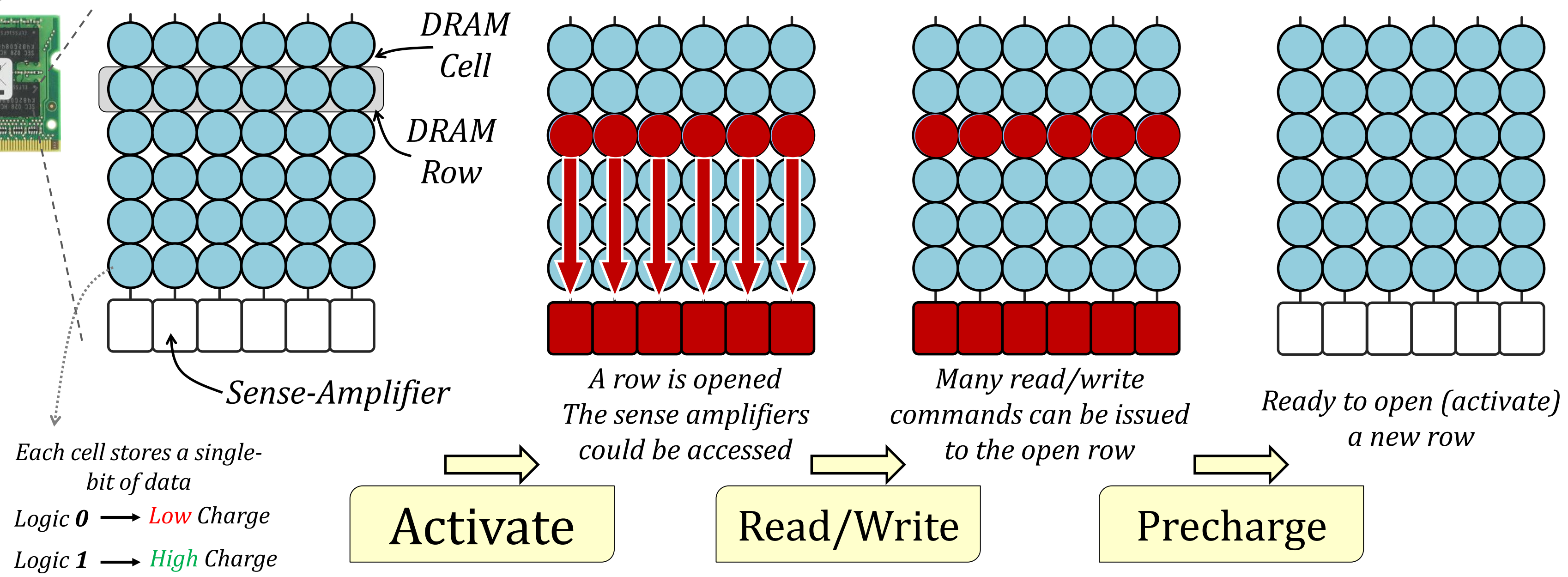
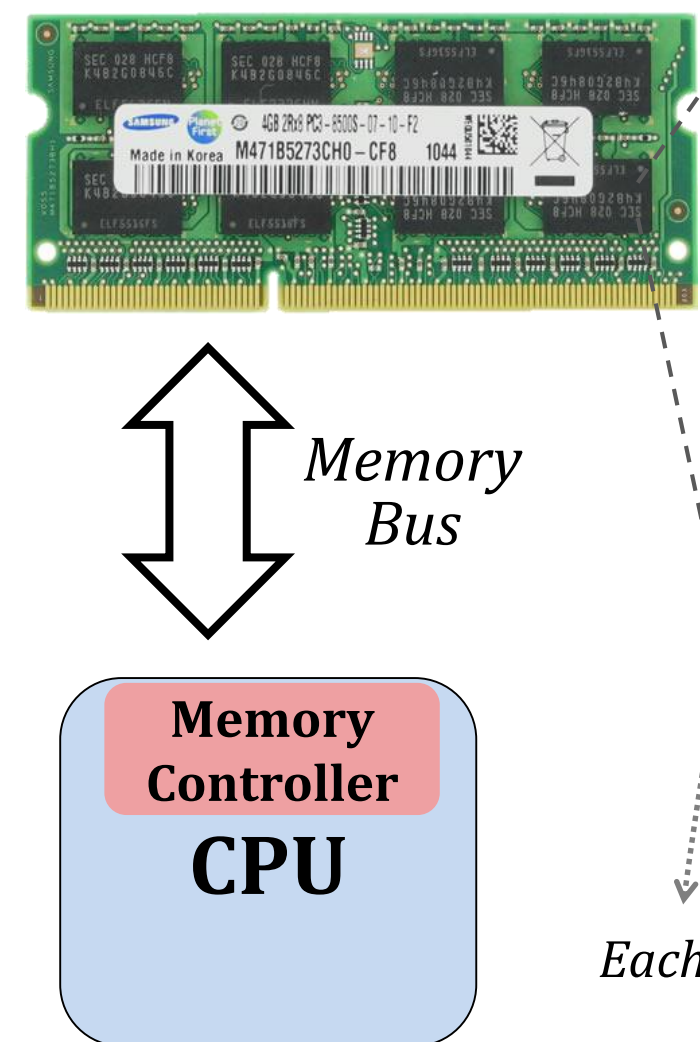


Use cases:

- **CROW-cache & CROW-ref** (20% speedup and 22% less DRAM energy)
- Mitigating RowHammer
- We hope CROW enables many other use cases going forward

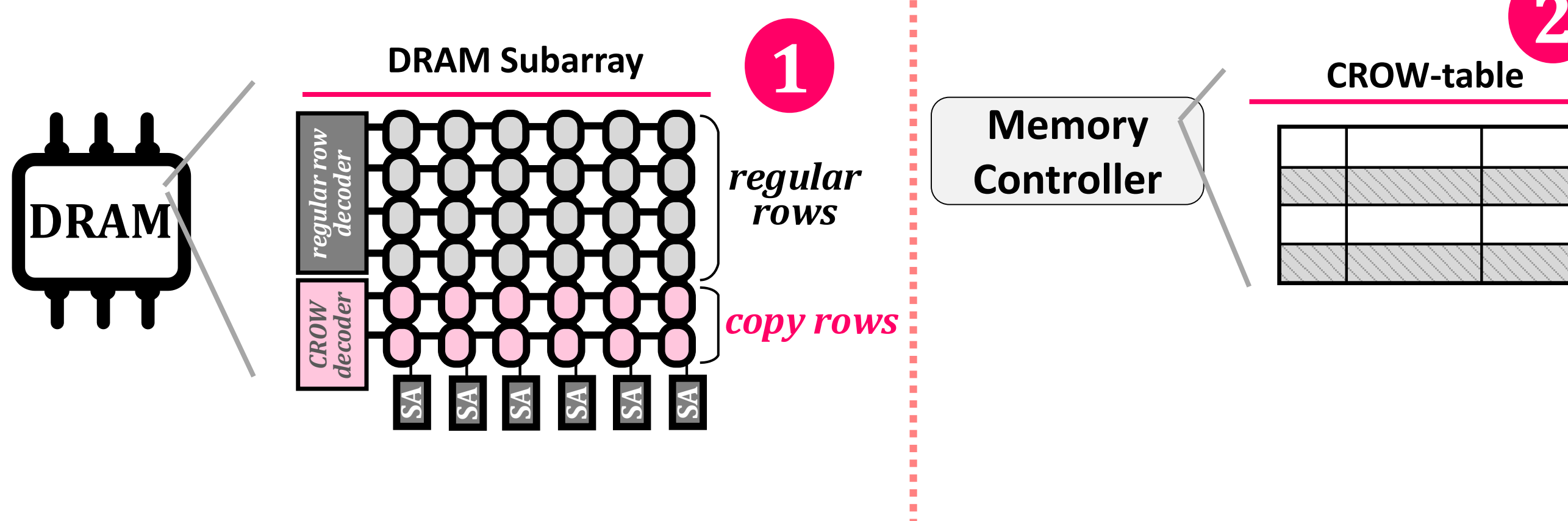
2: DRAM Operation Basics

DRAM Module



3: The Components of CROW

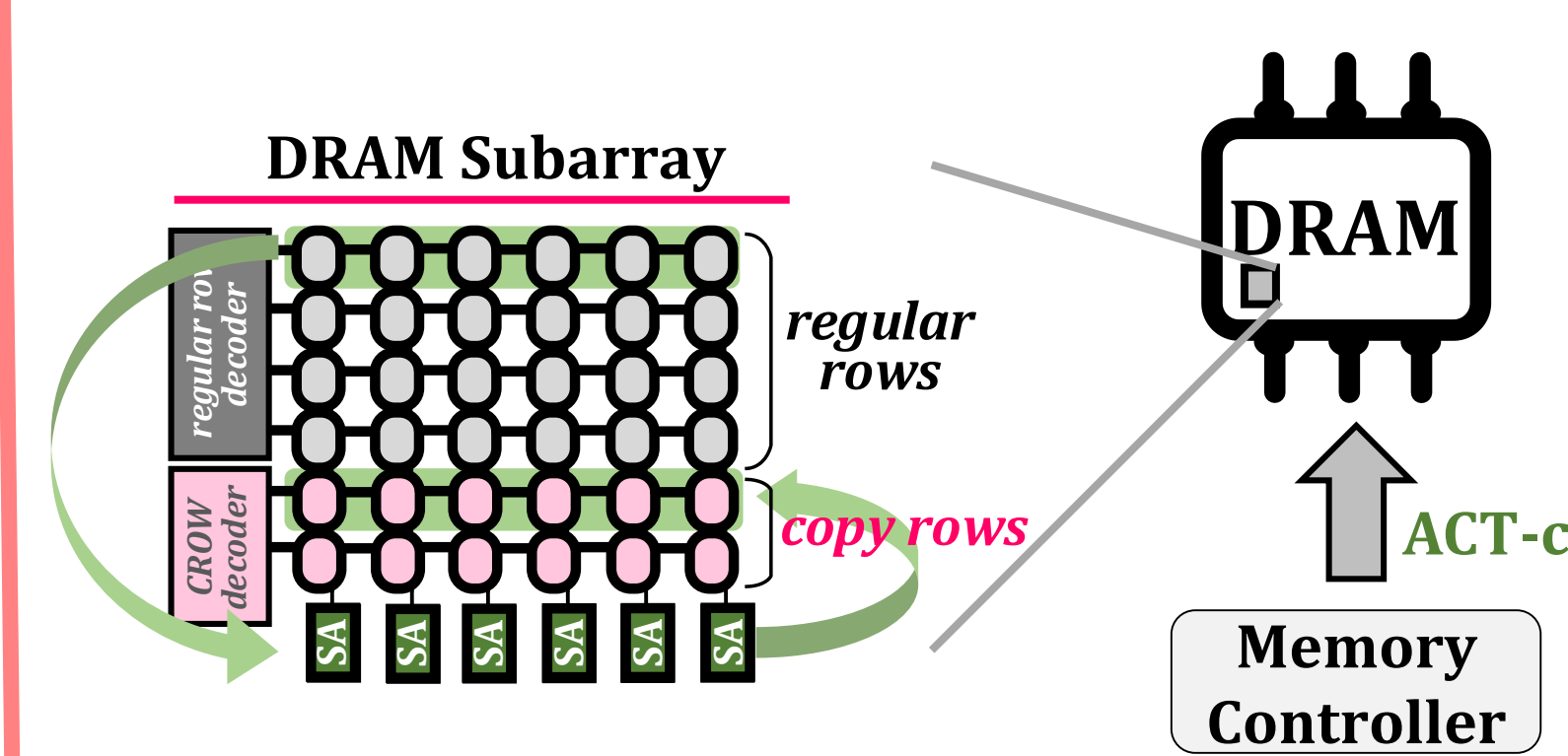
Copy-Row DRAM (CROW): a flexible in-DRAM substrate that can be used in multiple different ways to address the performance, energy efficiency, and reliability challenges of DRAM



4: CROW Operations

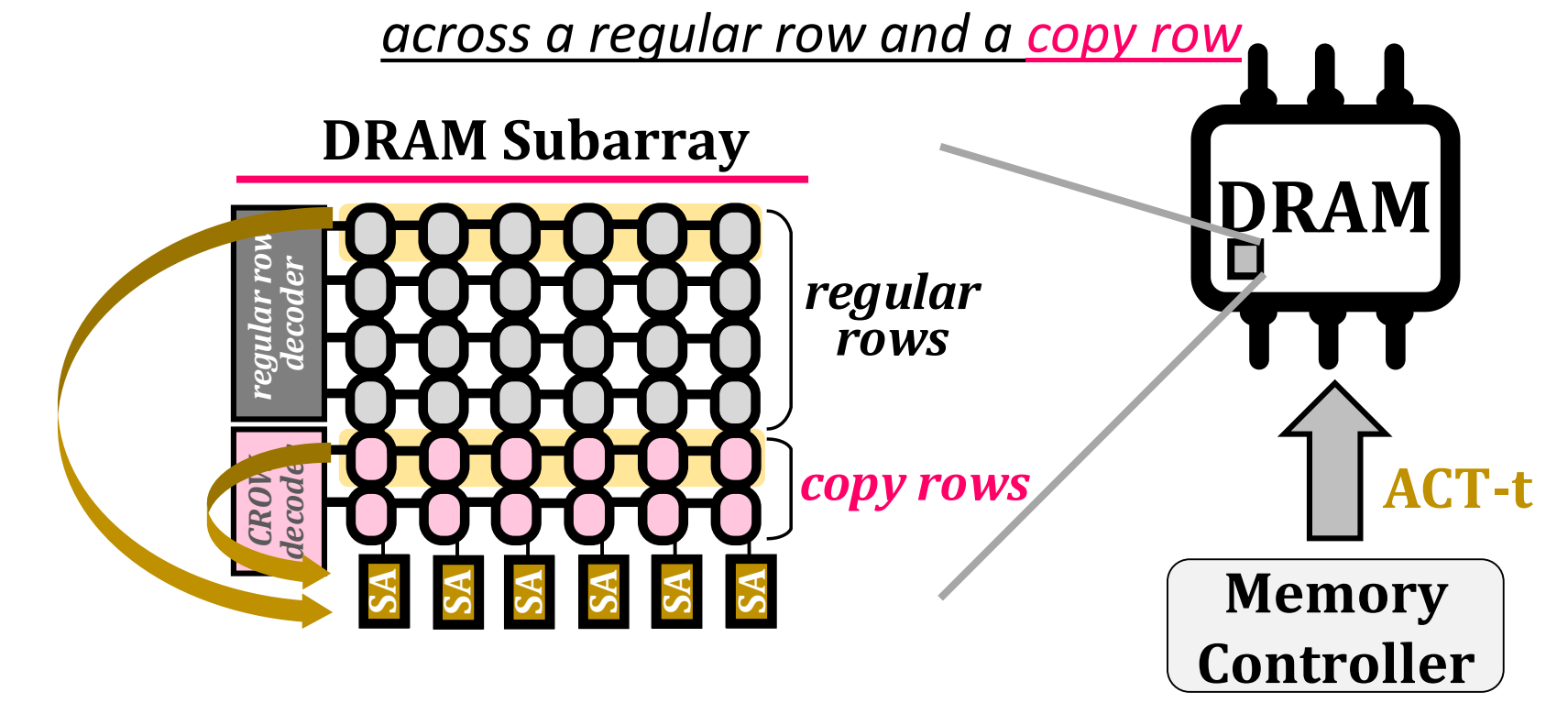
Row Copy

Enables quickly copying a regular row into a **copy row**



Two-Row Activation

Enables fast access to data that is duplicated across a regular row and a **copy row**



5: CROW-cache

Problem: High access latency

Key idea: Use **copy rows** to enable low-latency access to most-recently-activated regular rows in a subarray

CROW-cache combines:

- **row copy** → copy a newly activated regular row to a **copy row**
- **Two-row activation** → activate the regular row and **copy row** together on next access

Reduces activation latency by **38%**

6: CROW-ref

Problem: Refresh has high overheads. Weak rows lead to high refresh rate

- **weak row**: at least one of the row's cells cannot retain data correctly when refresh interval is increased

Key idea: Avoid storing data in a **weak** regular row by remapping it to a **strong copy row**

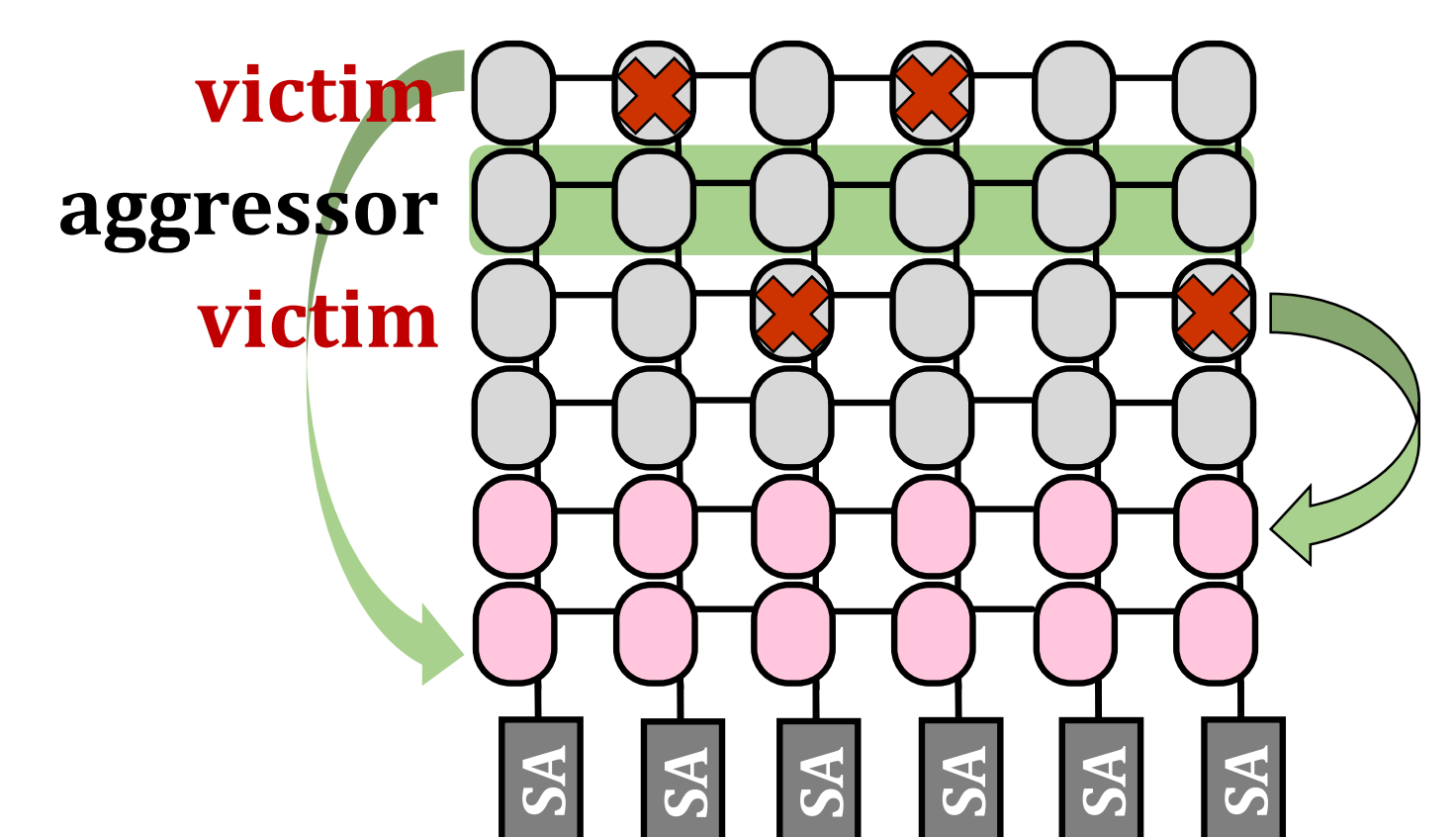
CROW-ref uses:

- **row copy** → copy a weak regular row to a strong **copy row**

CROW-ref eliminates more than half of the refresh requests

7: Mitigating RowHammer

Key idea: remap victim rows to copy rows

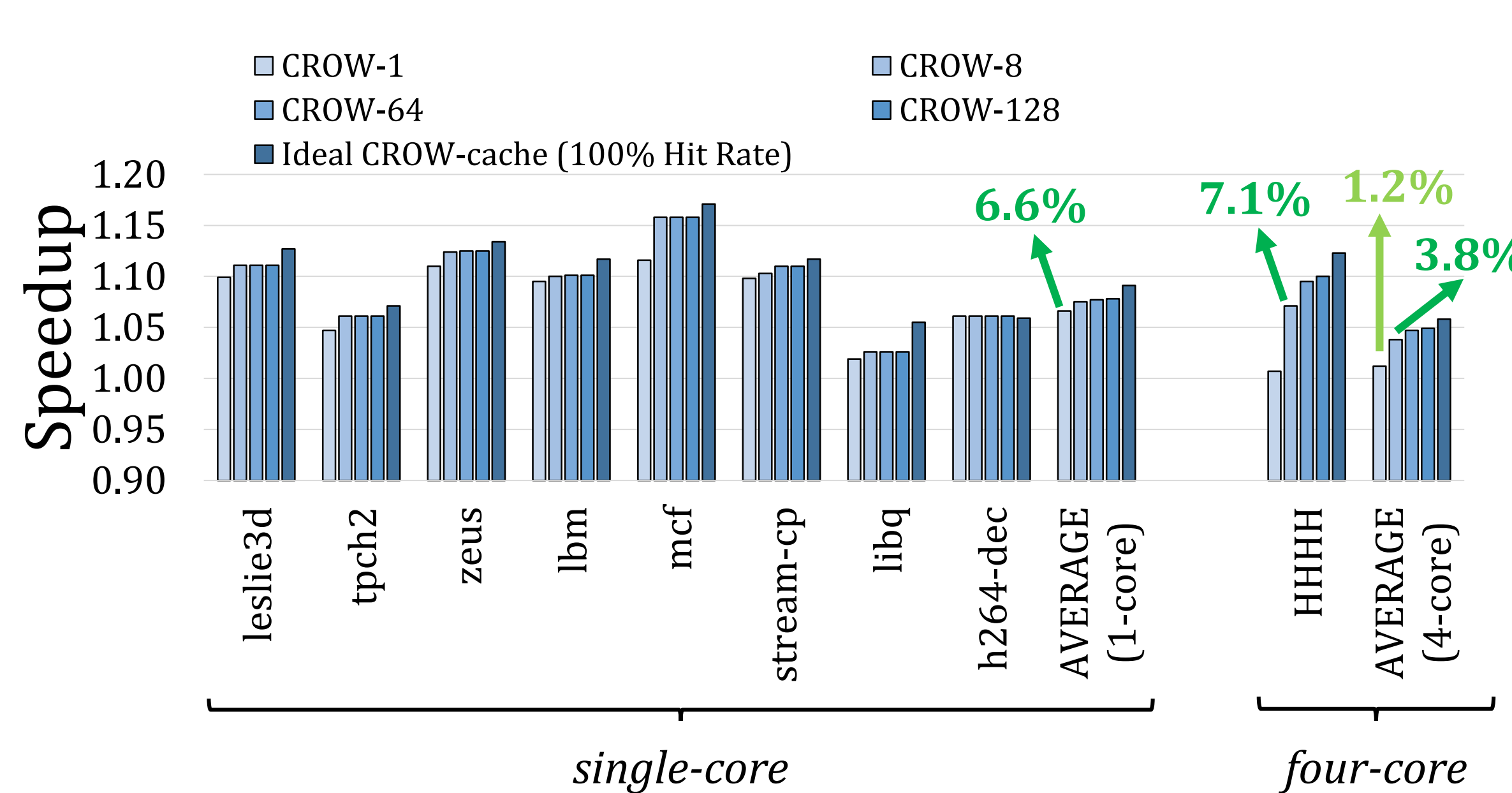


8: Evaluation

Methodology

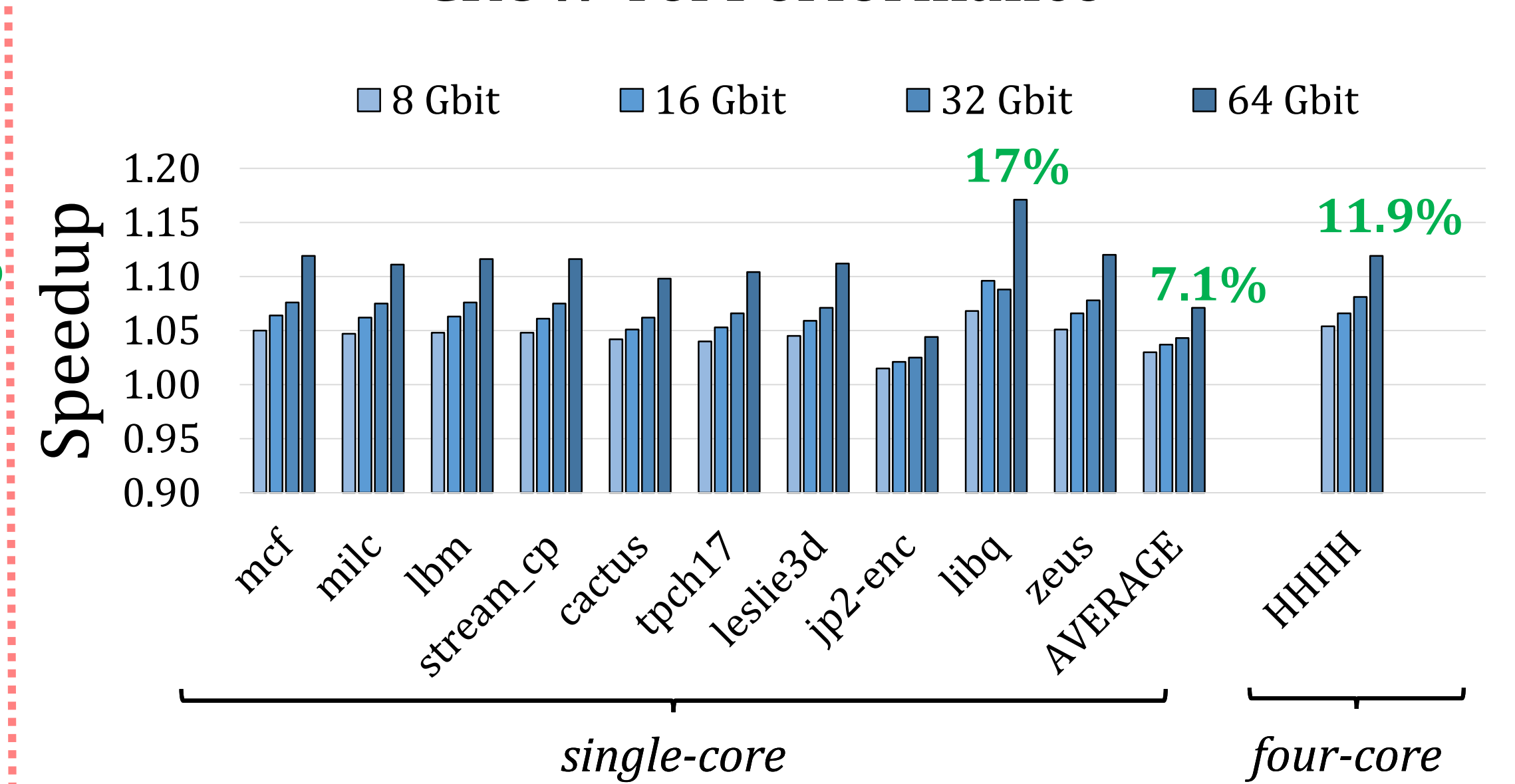
- **Simulator**
 - DRAM Simulator (Ramulator [Kim+, CAL'15])
<https://github.com/CMU-SAFARI/ramulator>
- **Workloads**
 - 44 single-core workloads
 - SPEC CPU2006, TPC, STREAM, MediaBench
 - 160 multi-programmed four-core workloads
 - By randomly choosing from single-core workloads
 - Execute at least 200 million representative instructions per core
- **System Parameters**
 - 1/4 core system with 8 MiB LLC
 - LPDDR4 main memory
 - 8 **copy rows** per 512-row subarray

CROW-cache Performance



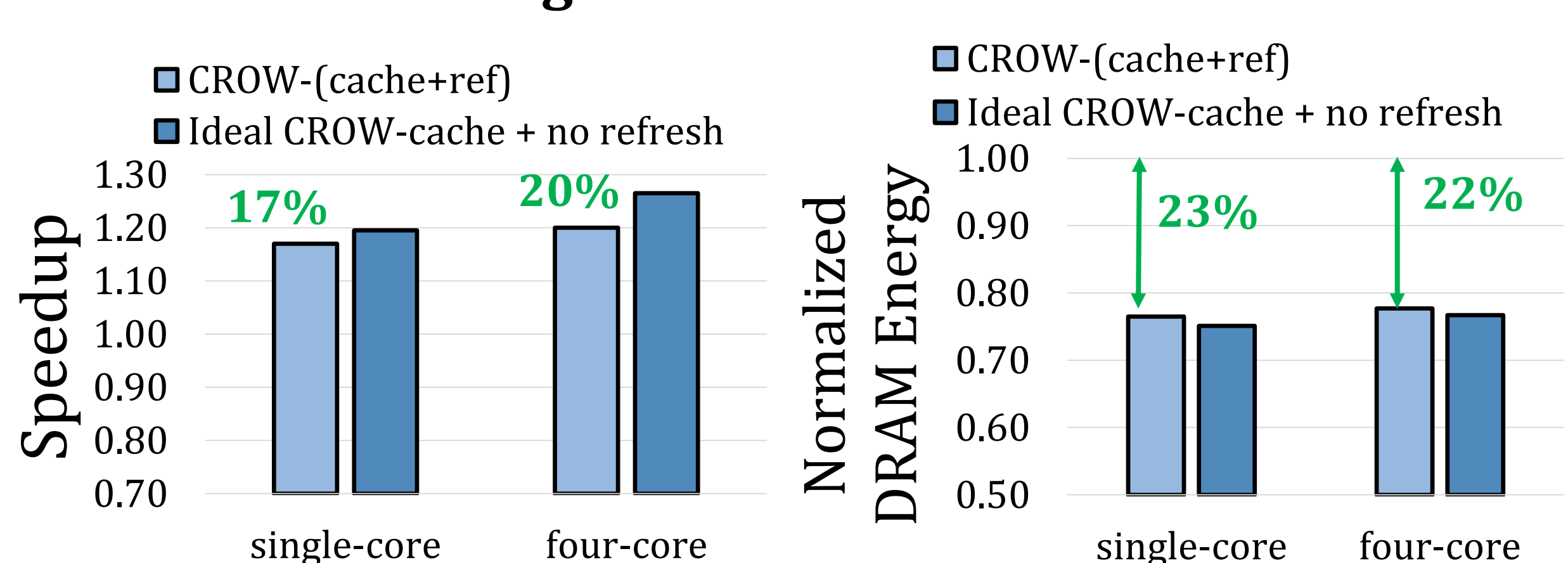
CROW-cache **improves** single-/four-core system performance

CROW-ref Performance



CROW-ref **significantly reduces** the performance overhead of DRAM refresh

Combining CROW-cache and CROW-ref



CROW-(cache+ref) provides more performance and DRAM energy benefits than each mechanism alone

Hardware Overhead

- For 8 copy rows:
- 0.5% DRAM chip area
 - 1.6% DRAM capacity
 - 11.3 KiB memory controller storage

Other Results in the Paper

- Sensitivity to:
 - Number of copy-rows per subarray
 - Chip density
 - Last-level cache capacity
- CROW-cache with prefetching
- CROW-cache against other in-DRAM caching mechanisms:
 - TL-DRAM [Lee+, HPCA'13]
 - SALP [Kim+, ISCA'12]

Available in July:

github.com/CMU-SAFARI/CROW

SAFARI