CS 211: Intro to Computer Architecture 3.1: Character and Integer Representations Cntd.

Minesh Patel Spring 2025 – Thursday 4 February

Announcements

• Due dates

- PA1 on Thursday (online via Gradescope)
- WA1 due Friday (online via Gradescope)

• WA2 and PA2 to be assigned on Friday (maybe Saturday...)

- WA2: 1 week
- PA2: 1.5 weeks

TA's recitation activities

• A variety of activities if you need more practice material or individual attention



Endianness

- •Two's Complement •Useful properties
- •Other signed representations
 - •Bias-K Representation
 - •Sign Magnitude
 - One's Complement

Webpage Encodings by Popularity

The following character encodings are used by less than 0.1% of the websites

UTF-8	98.5%
ISO-8859-1	1.1%
Vindows-1252	0.3%
Vindows-1251	0.2%
EUC-JP	0.1%
EUC-KR	0.1%
GB2312	0.1%
	W3Techs.com, 3 February 202

- Windows-1250
- Shift JIS
- ISO-8859-2
- Big5
- ISO-8859-15
- ISO-8859-9
- GBK
- US-ASCII
- Windows-1254Windows-874
- Windows-1256
- Windows-1255



- Windows-1253
- GB18030
- Windows-1257
- KOI8-R
- KS C 5601ISO-2022-JP
- ISO-2022-JP
 ISO-8859-4
- UTF-7
- ISO-8859-8
- ISO-8859-5
- ISO-8859-6
- Windows-31J
- ANSI_X3.110-1983KOI8-U
- Windows-1258
- ISO-8859-16
- ISO-8859-13
- Big5 HKSCS
- ISO-8859-3
- ISO-8859-10
- ISO-8859-14ISO-8859-11
- ISO-8859
 IBM850
- Windows-949

UTF-8 and UTF-16





Byte Order Ambiguity

- **Recall:** memory is a contiguous array of **bytes**
- Need a convention for how N-byte "words" are represented as bytes



Endianness More Generally

- Big and little endianness is just a matter of convention
 - Hardware and software must agree on that convention



• Endianness must be clearly specified for any **multi-byte** representation

Endianness in Action

• File containing 32-bit (N=4-byte) integers

mp2099@ilab3:~/cs211/experiment\$ file integers.dat integers.dat: data mp2099@ilab3:~/cs211/experiment\$ cat integers.dat; echo "3DUfw0000000

mp2099@ilab3:~/cs211/experiment\$ xxd -g 4 integers.dat 000000000: 00112233 44556677 8899aabb ccddeeff ..."3DUfw...... mp2099@ilab3:~/cs211/experiment\$ xxd -e integers.dat 00000000: 33221100 77665544 bbaa9988 ffeeddcc ..."3DUfw.....



•Endianness

•Two's Complement

Useful properties

•Other signed representations

- •Bias-K Representation
- •Sign Magnitude
- •One's Complement

Recap: Signed vs. Unsigned Numbers

• **Sign:** the plus/minus sign you put in front of a number



TIL Java doesn't have these 🛞

Unsigned Nonnegative integers

[0, inf) 0 10 31

When you're:

- Not doing arithmetic
 - ID numbers (no add/subtract)
 - Using bitwise operations
- 100% sure you don't need a negative number
 - Degrees Kelvin [o, inf)

Range

Examples

Use Cases

Signed

Positive or negative integers

(-inf, inf)

-10 10

Most use cases, really 😊

- Degrees Centigrade [-273, infinity]
- Integer arithmetic

Recap: Unsigned Number Representations

• It's just the binary representation of our value



• With N bits...





How can we represent **negative numbers** using bits?



Ideal Properties of a Signed Representation

• Arithmetic should "make sense"



Let C : value \rightarrow representation We want that C(0) = 0

Arithmetic on Values

Programmer's Job

Programmer and hardware

should do the same thing

Arithmetic on Representations

Hardware's Job

 $C(v_0) + C(v_1) = C(v_2)$ $C(v_0) - C(v_1) = C(v_2)$ $C(v_0) * C(v_1) = C(v_2)$

"Sensible" Arithmetic

• Operations on **representations** should mirror operations on **values**



Unsigned Representation Properties



 Easy to convert from value to representation

•
$$R = (V)_2$$

$$\checkmark C(0) = 0$$

- \checkmark Sensible arithmetic (mod 2³)
 - Overflow is OK
 - Programmer and hardware do the same thing

K Negative numbers

Extending Unsigned to Signed

Key idea: generate the negative numbers through subtraction
Find the "additive inverse" of each positive representation

$$(001)_{2} + X = (000)_{2}$$

$$(001)_{2} + X = (1000)_{2} \quad (equivalent mod 2^{3})$$

$$X = (1000)_{2} - (001)_{2}$$

$$X = (100)_{2} \rightarrow (-1)_{0}$$

$$X = (10)_{2} \rightarrow (-1)_{0}$$

Extending Unsigned to Signed

- Key idea: generate the negative numbers through subtraction
 - Find the "additive inverse" of each positive representation



=1000

Two's Complement Representation



• "Sign bit" helps make hardware efficient



Two's Complement Representation



Converting to 2's complement

$$\int (V)_2 \qquad V \ge 0$$

$$R = \begin{cases} (~|V| + 1)_2 & V < 0 \\ 1. "Flip" the bits of (|V|)_2 \\ 2. Add one \end{cases}$$

• Converting from 2's complement

$$\left((R)_{10} \right) MSB = 0$$

$$I = \begin{cases} -(~|R| + 1)_{10} & MSB = 1 \end{cases}$$

- 2. Add one
- 3. Add the minus sign

N-Bit Two's Complement to Decimal 5-67 +9=(01001) $(+9)_{10} = (01001)_2$ 10110 61001)2 -> (x)10 $-9 = (10111)_2$ $1.7^{3} + 1.2^{\circ} = (9)_{10}$ $(R)_{10}$ MSB = 0 $V = \begin{cases} -(\sim |R| + 1)_{10} \\ 1. "Flip" the bits of R \\ 2. Add one \\ 3. Add the minus sign \end{cases}$ MSB = 1 $(-9)_{10} = (10111)_2$ 01000

 $(-1001)_{7}$

 $(-9)_{10}$

Decimal to N-Bit Two's Complement

$$\mathbf{R} = \begin{cases} (V)_{2} & V \ge 0 \\ (\sim |V| + 1)_{2} & V < 0 \\ 1. "Flip" the bits of (|V|)_{2} \\ 2. Add one \end{cases}$$

- Always double-check the sign bit
 - V must fit in N-1 bits

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 $(+9)_{10} = (01001)_2$ $(9)_{10} = (1001)_{7}$ 4 N=4bit= $(-9)_{10} = (10111)_2$ $(9)_{0}=(100)_{2}$ (0100) 0110 10110² 6111

Sanity Checks with Two's Complement

- Check the MSB: sign is positive (0) or negative (1)
- Check the LSB: the number is even (0) or odd (1)

$$(+8)_{10} = (01000)_2$$

$$(+9)_{10} = (01001)_2$$

$$(-8)_{10} = (11000)_2$$

 $(-9)_{10} = (10111)_2$

Two's Complement Properties



X Strange conversion from value to representation

•
$$R = \begin{cases} (V)_2 & V \ge 0 \\ (\sim |V| + 1)_2 & V < 0 \end{cases}$$

$$\checkmark C(0) = 0$$

- \checkmark Sensible arithmetic (mod 2³)
 - Overflow is OK
 - Programmer and hardware do the same thing
- Negative numbers

Integer Representations

Unsigned

- \checkmark Simple code
- $\checkmark C(0) = 0$
- ✓ Sensible arithmetic
- X Negative numbers

Two's Complement

- X Simple code
- $\checkmark C(0) = 0$
- 🗸 Sensible arithmetic
 - \checkmark Negative numbers



•Endianness

•Two's Complement •Useful properties

•Other signed representations

- •Bias-K Representation
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Two's Complement and Positional Notation

• MSB's weight is now a negative power of two

$$(+9)_{10} = -0.2^{4} + 1.2^{3} + 0.2^{2} + 0.2^{1} + 1.2^{0}$$

$$(-9)_{10} = -1 \cdot 2^{4} + 0 \cdot 2^{3} + 1 \cdot 2^{2} + 1 \cdot 2^{1} + 1 \cdot 2^{0}$$

-16 $8 + 4 + 2 + 1 = 15 = -1$

Unsigned and Two's Complement Arithmetic



- Unsigned and Two's Complement arithmetic is identical by construction
 - We defined negative numbers to be equivalent to positive numbers mod 2³
 - You can add/sub/mul two's complement representations as if they are unsigned
 - You'll always get the right answer (mod 2³)

Unsigned and Two's Complement in C

Unsigned

- unsigned char
- unsigned short
- unsigned int
- unsigned long
- unsigned long long
- uint{8,16,32,64}_t







•Endianness

•Two's Complement •Useful properties

Other signed representations

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Integer Representations

Most computers use these

Unsigned

✓ Simple code

$$\checkmark C(0) = 0$$

- ✓ Sensible arithmetic
- X Negative numbers

Two's Complement

X Simple code

$$\checkmark C(0) = 0$$

- ✓ Sensible arithmetic
- ✓ Negative numbers

Bias-K

Sign Magnitude

Part of float representation (also used in most computers)

One's Complement

Sometimes useful



•Endianness

- •Two's Complement •Useful properties
- •Other signed representations
 - •Bias-K Representation
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Bias-K Representation

• Key idea: shift all values by some integer K



K = 10







32

Bias-K Arithmetic



Arithmetic is weird for K≠0

Arithmetic with **one** value is OK (0-K) + 1 = (1-K)000 + 001 = 001

Arithmetic with **two** values is NOT $(0-K) + (1-K) = (1-2K) \times 000 + 001 = 001$

Bias-K Representation



- Easy to convert from value to representation
 - $R = (V-K)_2$
- (0) = 0

 \times Sensible arithmetic (mod 2³)

- Add/subtract are weird
- Might not even have a zero

✓ Negative numbers

Bias-K with "Standard Bias"

 $K = 2^{N-1}$



- Same values as 2C, different representation
 - Wheel is rotated 180 degrees
 - MSB is inverted
- Bias-K preserves total order over representations
 Useful for logical comparisons on the representations



$$Largest Number$$

$$1 1111...1111$$

$$Value = 2^{N} - 1 - K$$

Integer Representations

Unsigned

✓ Simple code

$$\checkmark C(0) = 0$$

- ✓ Sensible arithmetic
- X Negative numbers

Two's Complement

X Simple code

$$\checkmark C(0) = 0$$

- Sensible arithmetic
- ✓ Negative numbers

Bias-K

Sign Magnitude

One's Complement

- ✓ Simple code
- X C(0) = 0
- X Sensible arithmetic
- ✓ Negative numbers



•Endianness

- •Two's Complement •Useful properties
- Other signed representations
 Bias-K Representation
 - •Sign Magnitude
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Sign Magnitude Representation

• Key idea: MSB indicates positive (0) or negative (1)



Sign Magnitude Representation

• Comparable to human notation (e.g., -27 vs +27)



Largest Number

$$\begin{array}{c|c}
0 \\
1 \\
\hline Value = +(2^{N-1}-1) \\
\end{array}$$

Sign Magnitude Arithmetic



Sign Magnitude Properties



- Easy to convert from value to representation
 - $R = [sign | (|V|)_2]$

$$\checkmark C(0) = 0$$

- X Sensible arithmetic (mod 2³)
 - Two zeroes
 - Values are non-monotonic
 - Add/subtract don't work

✓ Negative numbers

Integer Representations

Unsigned

✓ Simple code

$$\checkmark C(0) = 0$$

- ✓ Sensible arithmetic
- × Negative numbers

Two's Complement

X Simple code

$$\checkmark C(0) = 0$$

- ✓ Sensible arithmetic
- ✓ Negative numbers

Bias-K

- ✓ Simple code
- (0) = 0
- X Sensible arithmetic
- ✓ Negative numbers

Sign Magnitude

- \checkmark Simple code
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One's Complement



•Endianness

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One's Complement Representation

•Key idea: correct the total order on sign magnitude representation



One's Complement Representation



One's Complement Properties



X Strange conversion from
value to representation $R = \begin{cases} 0 & | & (V)_2 & V \ge 0 \\ 1 & | & (\sim |V|)_2 & V < 0 \end{cases}$

$$\checkmark C(0) = 0$$

- X Sensible arithmetic (mod 2³)
 - Two zeroes
 - Add/subtract make no sense

✓ Negative numbers

Integer Representations

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Two's Complement

X Simple code

$$\checkmark C(0) = 0$$

- Sensible arithmetic
- ✓ Negative numbers

Bias-K

- ✓ Simple code
- (0) = 0
- X Sensible arithmetic
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Sign Magnitude

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One's Complement

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Thursday: Floating-Point



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Spring 2025 – Thursday 30 January