# **ACM SIGMOD Programming Contest 2023**

**TEAM X2A3008M** 



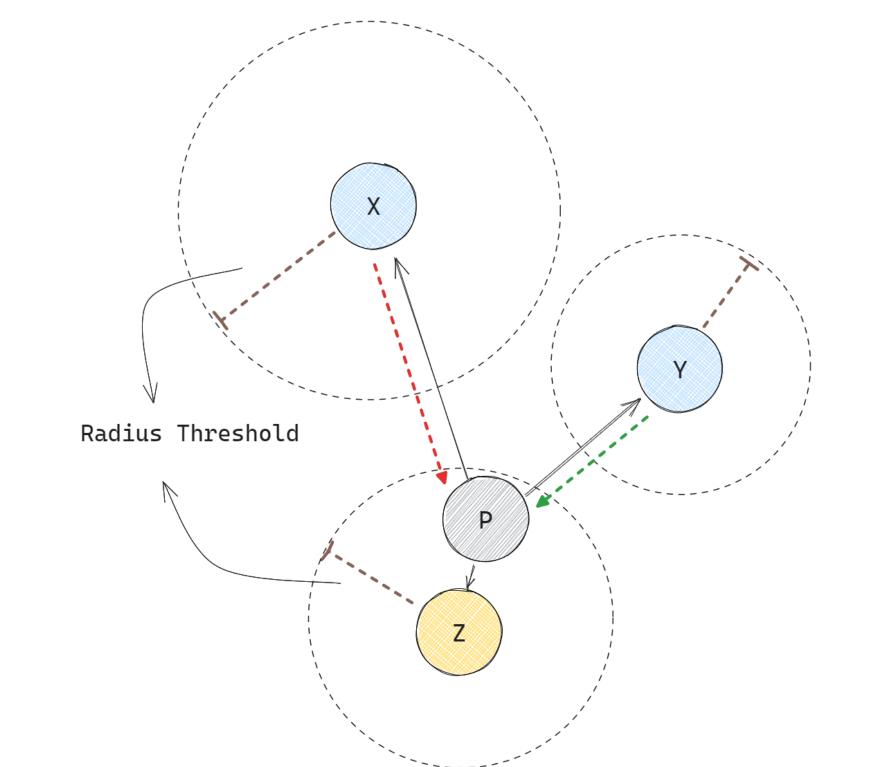
Member: Meng Chen, Advisor: Kai Zhang DASlab - System Group @ Fudan University

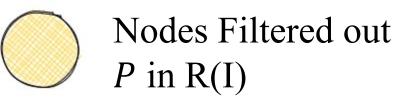


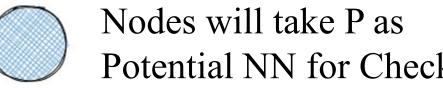
Contact: mengchen22@m.fudan.edu.cn, zhangk@fudan.edu.cn

## **Task Overview**

**Task:** The objective is to create an approximate K-NN graph using a set of high-dimensional vectors where each vertex is linked to its approximate k nearest neighbors based on the Euclidean distance. **Dataset:** 







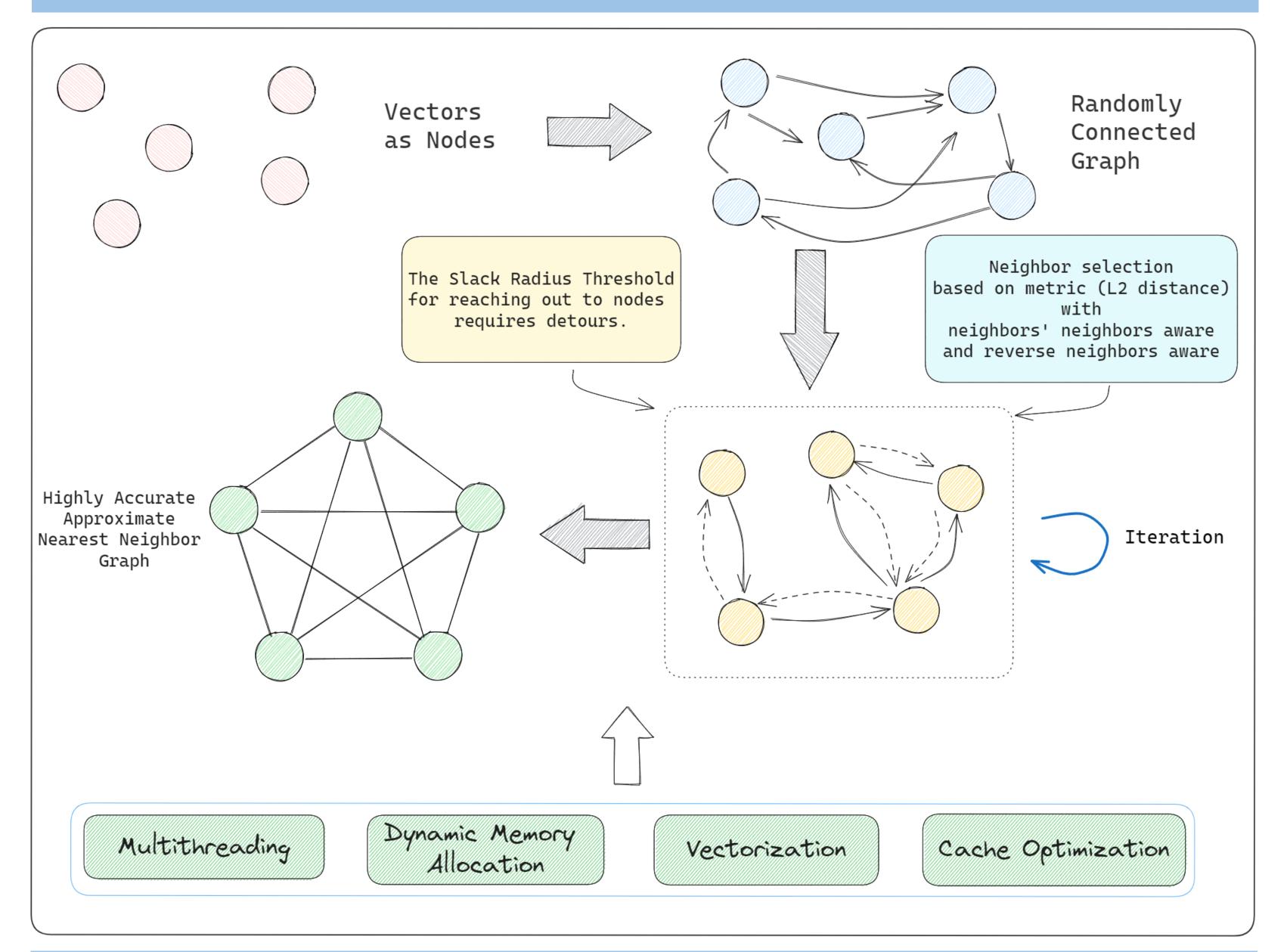
Potential NN for Checking

Dataset	# of vectors	dimension
Turing	10,000,000	100

Vectors are Bing queries encoded by Turing AGI v5.

**Output:** Approximate KNN Graph (K=100) **Input:** A set of vectors **Evaluation Metric:**  $Recall = \frac{ground \ truth (100-NN \ neighbors)}$ 100 Hardware Conditions: Azure Standard F32s v2, 32C64G **Time Limit:** 30 minutes + 60 seconds (reprozip overhead)

## **Solution Overview**



Fail to Establish Links due to **Reverse-NN Set Truncated** 

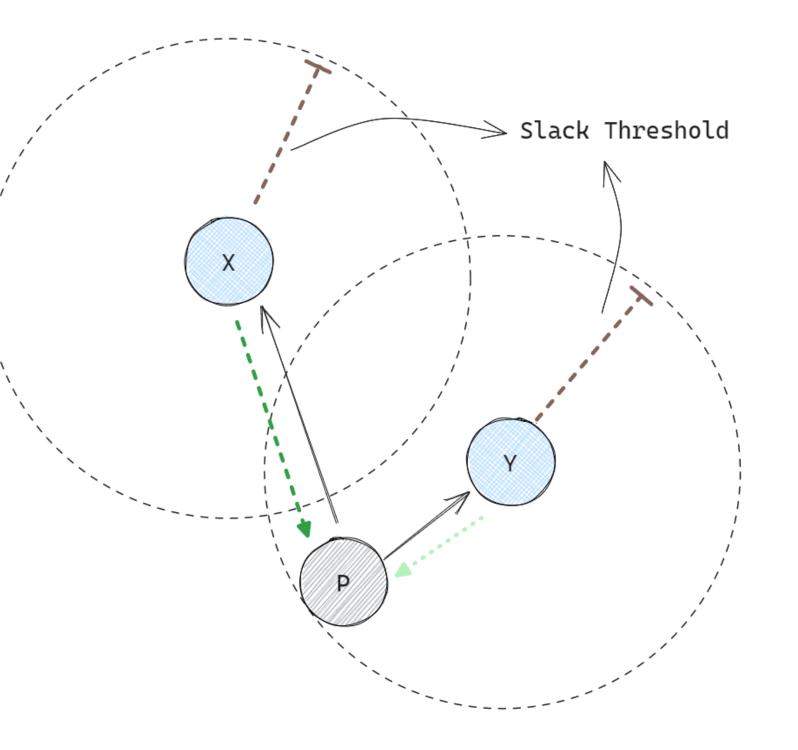
**Goal**: preserve  $Y \rightarrow P$  and add  $X \rightarrow P$ 

Unfortunately, the curse of dimensionality frequently causes hubs to appear in high-dimensional space, which we also observed in the Turing Dataset during our analysis. This causes the reverse-NN set to be truncated, thereby saving memory budget, but also undermining neighbor discovery through reverse-NN paths.

#### With Slack-Threshold:

This method is based on two key principles:

- A neighbor of a neighbor is also likely to be a neighbor. (from NN-Descent)
- A closer neighbor could be found by taking a slight detour. (Inspired by the RNG property [2])



# Implementation

Algorithm: Slack-Threshold NN-Descent, which is based on NN-Descent [1] with optimizations and improvements.

**Data Structures Required by Each Vector(Node):** 

- Candidates Set (C)

With Slack-Threshold, nodes with relatively far distance from *P* are more likely to include *P* into its neighbor checking phase.(Node X)

It filters out Node Y. However, by applying local join and relying on principle 1 of NN-Descent, we ensure that local close nodes are detected and well-connected. This approach increases the likelihood of reaching nodes that would otherwise require detours, while also preserving strong local connections.

#### **Low-Level Design:**

- The previously mentioned structure will be dynamically maintained to reduce memory footprint.
- SIMD with AVX-512 is used with aligned data for accelerating L2 distance computation.
- Prefetching has been effectively integrated into the computational checking progress.

- Incoming Nodes (I)
- Reverse-NN Nodes (R)

#### **Overall Procedure:**

- Randomly selecting 100 neighbors for each node.
- The incoming nodes are recognized as closer NNs in the current round (Random nodes at the beginning). The reverse-NN nodes consist of nodes that have out-degrees to the current one.
- 3. Pairs generated by Cartesian Product between set  $I \cup R(I)$  itself, set  $I \cup R(I)$  and set  $C \cup R(C)$  for each node are mutually checked whether they constitute the nearest neighbor.

#### With Tight-Threshold:

The original algorithm utilizes the radius threshold to establish the range for reverse-NN detection and also filter intimate nodes, which are considered well-connected due to their shorter distances.

#### Results are selected from submissions in the contest.

Dataset	Recall	Time (s)
Turing-10M	0.981	1847
Turing-10M	0.976	1654
Turing-10M	0.954	1300

### References

[1] Dong, Wei, Charikar Moses, and Kai Li. "Efficient k-nearest neighbor graph construction for generic similarity measures." Proceedings of the 20th international conference on World wide web. 2011.

[2] Jaromczyk, Jerzy W., and Godfried T. Toussaint. "Relative neighborhood graphs" and their relatives." Proceedings of the IEEE 80.9 (1992): 1502-1517.