

Internet Technology
02. Network Protocol Layers & Sockets

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Protocols

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What's in the data?

- For effective communication
 - same language, same conventions
- For computers:
 - electrical encoding of data
 - where is the start of the packet?
 - which bits contain the length?
 - is there a checksum? where is it? how is it computed?
 - what is the format of an address?
 - byte ordering

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Protocols

These instructions & conventions are known as **protocols**

Protocols encompass data formats, order of messages, responses

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Layering

To ease software development and maximize flexibility:

- Network protocols are generally organized in **layers**
- Replace one layer without replacing surrounding layers
- Higher-level software does not have to know how to format an Ethernet packet
 - ... or even know that Ethernet is being used

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Protocols

Exist at different levels

understand format of address and how to compute a checksum *humans vs. whales*
different wavelengths

versus

request web page *French vs. Hungarian*

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Layering

Most popular model of guiding (not specifying) protocol layers is

OSI reference model

Adopted and created by ISO

7 layers of protocols

OSI = Open Systems Interconnection
From the ISO = International Organization for Standardization

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OSI Reference Model: Layer 1

Transmits and receives raw data to communication medium

Does not care about contents

Media, voltage levels, speed, connectors

Deals with representing bits

1

Physical

Examples: USB, Bluetooth, 802.11

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OSI Reference Model: Layer 2

Detects and corrects errors

Organizes data into **frames** before passing it down. Sequences packets (if necessary)

Accepts acknowledgements from immediate receiver

Deals with frames

2

Data Link

1

Physical

Examples: Ethernet MAC, PPP

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OSI Reference Model: Layer 3

Relay and route information to destination

Manage journey of **datagrams** and figure out intermediate hops (if needed)

Deals with datagrams

3

Network

2

Data Link

1

Physical

Examples: IP, X.25

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OSI Reference Model: Layer 4

Provides an interface for end-to-end (application-to-application) communication: sends & receives **segments** of data. Manages flow control. May include end-to-end reliability

Network interface is similar to a mailbox

Deals with segments

4

Transport

3

Network

2

Data Link

1

Physical

Examples: TCP, UDP

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OSI Reference Model: Layer 5

Services to coordinate dialogue and manage data exchange

Software implemented switch

Manage multiple logical connections

Keep track of who is talking: establish & end communications

Deals with data streams

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Session

4

Transport

3

Network

2

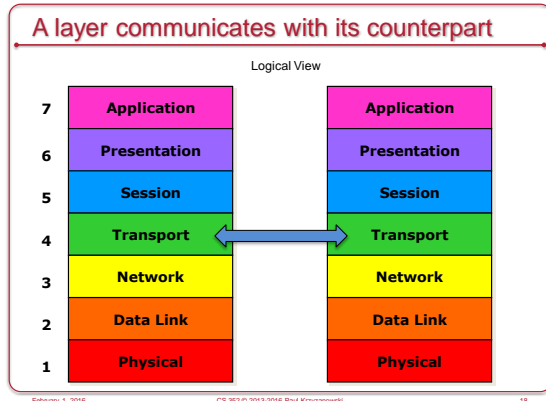
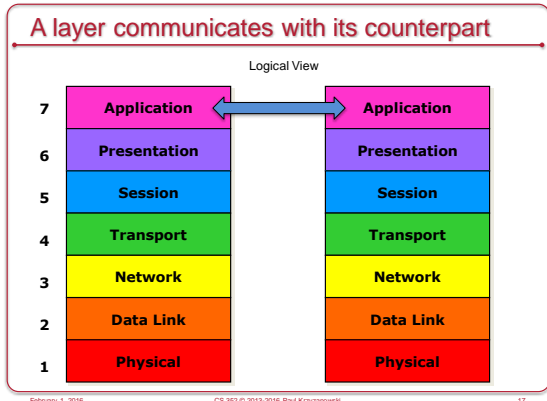
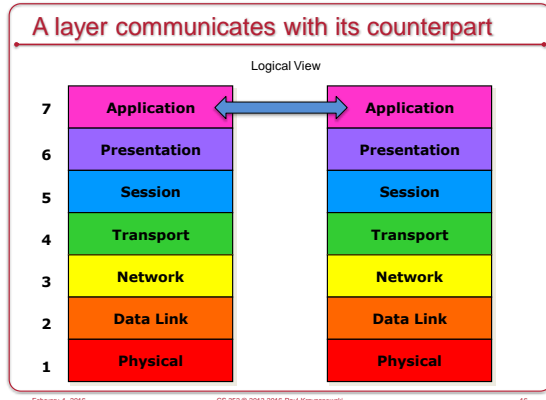
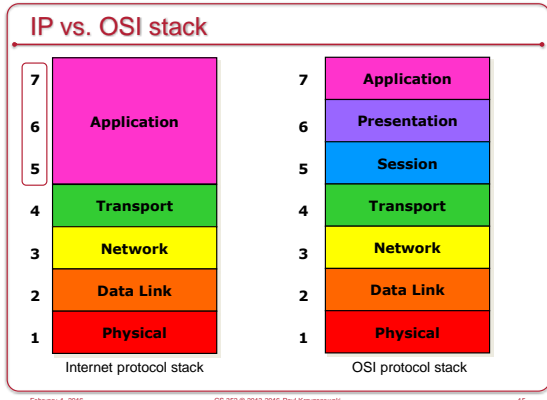
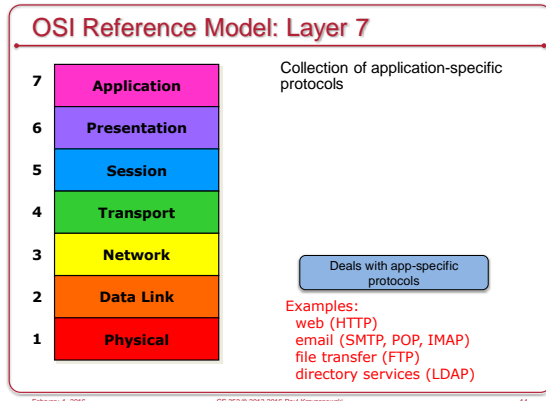
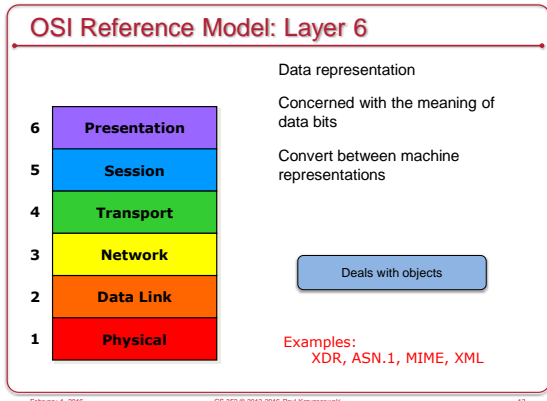
Data Link

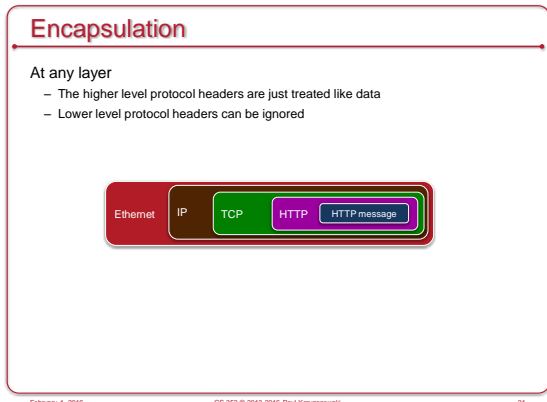
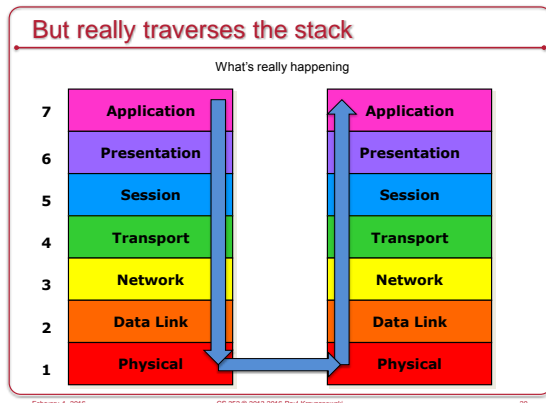
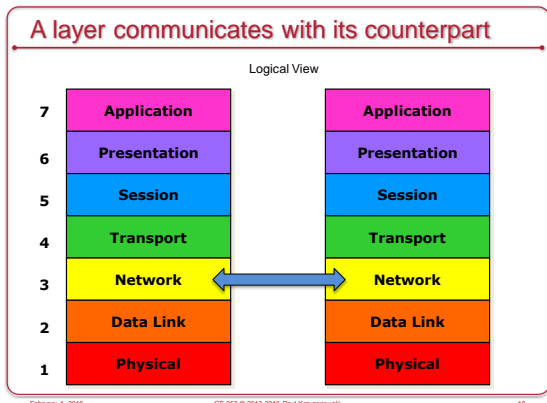
1

Physical

Examples: HTTP 1.1, SSL

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The Application Layer

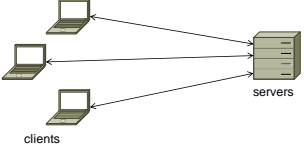
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- ### Writing network applications
- Network applications communicate with each other over a network
- Regular processes running on computers
 - Any process can access the network
 - Use a *network API* to communicate
 - The app developer does not have to program the lower layers
 - Speak a well-defined *application-layer protocol*
 - If the protocol is well-defined, the implementation language does not matter
 - E.g., Java on one side, C on the other
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- ### Application Architectures
- Client-server
 - Peer-to-peer (P2P)
 - Hybrid
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Client-Server architecture

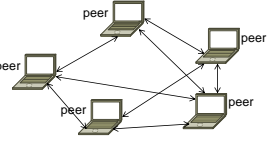
- Clients send requests to a server
- The server is always on and processes requests from clients
- Clients do not communicate with other clients
- Examples:
 - FTP, web, email



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Peer-to-Peer (P2P) architecture

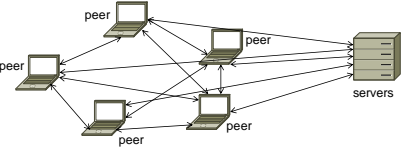
- Little or no reliance on servers
- One machine talks to another (**peers**)
- Peers are not owned by the service provider but by end users
- **Self-scalability**
 - System can process more workload as more machines join
- Examples
 - BitTorrent, Skype



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Hybrid architecture

- Many peer-to-peer architectures still rely on a server
 - Look up, track users
 - Track content
 - Coordinate access
- But traffic-intensive workloads are delegated to peers



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It's always (*mostly*) client-server!

Even for P2P architectures, we may use client-server terminology

- **Client**: process making a request
- **Server**: process fulfilling the request

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Network API

- App developers need access to the network
- A **Network Application Programming Interface (API)** provides this
 - Core services provided by the operating system
 - Operating System controls access to resources (the network)
 - Libraries handle the rest

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What do we need as programmers?

- **Reliable data transfer**
 - Reliable delivery of a stream of bytes from one machine to another
 - In-order message delivery
 - **Loss-tolerant applications**
 - Can handle unreliable data streams
- **Throughput**
 - **Bandwidth sensitive applications**: require a particular bitrate
 - **Elastic applications**: can adapt to available bitrate
- **Delay & Jitter Control**
 - Jitter = variation in delay
- **Security**
 - Authentication of endpoints, encryption of content, assured data integrity

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What IP gives us

IP give us two transport protocols

- **TCP: Transmission Control Protocol**
 - **Connection-oriented service**
 - Operating system keeps state
 - **Full-duplex connection**: both sides can send messages over the same link
 - **Reliable data transfer**: the protocol handles retransmission
 - **In-order data transfer**: the protocol keeps track of sequence numbers
- **UDP: User Datagram Protocol**
 - **Connectionless service**: lightweight transport layer over IP
 - Data may be lost
 - Data may arrive out of sequence
 - Checksum for corrupt data: operating system drops bad packets

What IP does not give us

- Throughput (bandwidth) control
- Delay and jitter control
- Security

We'll see how these were addressed later in the course

Usually addressed at the application with protocols such as SSL. Stay tuned for VPNs...

Addressing machines

(We'll examine IP addresses in depth later)

Machine addresses

- We identify machines with IP addresses: 32-bit numbers
- Example

cs.rutgers.edu = 128.6.4.2 = 0x80060402

Addressing applications

Communication endpoint at the machine

- **Port number**: 16-bit value
- Port number = transport endpoint
 - Allows application-application communication
 - Identifies a specific data stream
- Some services use well-known port numbers (0 – 1023)
 - IANA: Internet Assigned Numbers Authority (www.iana.org)
 - Also see the file /etc/services
- **ftp: 21/TCP ssh: 22/tcp smtp: 25/tcp http: 80/tcp ntp: 123/udp**
- Ports for proprietary apps: 1024 – 49151
- Dynamic/private ports: 49152 – 65535

The Application Layer: Sockets

Sockets

- Dominant API for transport layer connectivity
- Created at UC Berkeley for 4.2BSD Unix (1983)
- Design goals
 - Communication between processes should not depend on whether they are on the same machine
 - Communication should be efficient
 - Interface should be compatible with files
 - Support different protocols and naming conventions
 - *Sockets is not just for the Internet Protocol family*

What is a socket?

Abstract object from which messages are sent and received

- Looks like a file descriptor
- Application can select particular style of communication
 - Stream (connection-oriented), datagram (connectionless), message-based, in-order delivery
- Unrelated processes should be able to locate communication endpoints
 - Sockets can have a name
 - Name should be meaningful in the communications domain
 - E.g., Address & port for IP communications

How are sockets used?

Client: web browser Server: web server

Send HTTP request message to get a page → Receive HTTP request message

Process HTTP request

Send HTTP response message → Receive HTTP response message

Display a page

Connection-Oriented (TCP) socket operations

Connectionless (UDP) socket operations

The sockets system call interface

POSIX system call interface

System call	Function
socket	Create a socket
bind	Associate an address with a socket
listen	Set the socket to listen for connections
accept	Wait for incoming connections
connect	Connect to a socket on the server
read/write, sendto/rcvfrom, sendmsg/rcvmsg	Exchange data
close/shutdown	Close the connection

Step 1 (client & server)

Create a socket

```
int s = socket(domain, type, protocol)
```

New socket

AF_INET
Address Family: group of protocols for communication.
AF_INET is for IPv4
AF_INET6 is IPv6
AF_BTH is Bluetooth

SOCK_STREAM
Type of protocol within the family
SOCK_STREAM: reliable, in-order, 2-way, TCP/IP
SOCK_DGRAM: datagrams (UDP/IP)
SOCK_RAW: "raw" - allows app to modify the network layer header

useful if some families have more than one protocol to support a given service.
0: unspecified

Conceptually similar to `open` BUT
- `open` creates a new reference to a possibly existing object
- `socket` creates a new instance of an object

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Step 2 (client & server)

Name the socket (assign address, port)

```
int error = bind(s, addr, addrlen)
```

socket
The socket from the socket system call.

Address structure
`struct sockaddr*`
This is a data structure that makes sense for whatever address family you selected.

length of address structure

Naming for an IP socket is the process of assigning our address to the socket. The address is the full transport address: the IP address of the network interface as well as the UDP or TCP port number

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Step 3a (server)

Set socket to be able to accept connections

```
int error = listen(s, backlog)
```

socket
The socket from the socket system call.

queue length for pending connections
Number of connections you'll allow between accept system calls

The socket that the server created with `socket` is now configured to accept new connections. This socket will *only* be used for accepting connections. Data will flow onto another socket.

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Step 3b (server)

Wait for a connection from client

```
int snw = accept(s, clntaddr, &clntalen)
```

socket
This is the listening socket

pointer to address structure
This tells you where the socket came from: full transport address.

length of address structure

new socket for this communication session

Block the process until an incoming connection comes in.

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Step 3 (client)

Connect to server

```
int error = connect(s, svraddr, svraddrlen)
```

socket
The socket from which we're connecting.

address structure
`struct sockaddr*`
Full transport address of the destination: address and port number of the service.

length of address structure

The client can send a connection request to the server once the server did a `listen` and is waiting for `accept`.

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Step 4. Exchange data

for connection-oriented service

- read/write system calls** (same as for file systems)
- send/rcv system calls**

```
int send(int s, void *msg, int len, uint flags);
int rcv(int s, void *buf, int len, uint flags);
```

Like read and write but these support extra flags, such as bypassing routing or processing out of band data. Not all sockets support these.
- sendto/rcvfrom system calls**

```
int sendto(int s, void *msg, int len, uint flags,
           struct sockaddr *to, int tolen);
int rcvfrom(int s, void *buf, int len, uint flags,
            struct sockaddr *from, int *fromlen)
```

If we're using UDP (connectionless), we don't need to do `connect`, `listen`, `accept`. These calls allow you to specify the destination address (`sendto`, `sendmsg`) to send a message and get the source address (`rcvfrom`, `rcvmsg`) when receiving a message.
- sendmsg/rcvmsg system calls**

```
int sendmsg(int s, struct msghdr *msg, uint flags);
int rcvmsg(int s, struct msghdr *msg, uint flags);
```

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Step 5

Close connection

```
shutdown(s, how)
```

how:
 SHUT_RD (0): can send but not receive
 SHUT_WR (1): cannot send more data
 SHUT_RDWR (2): cannot send or receive (=0+1)

You can use the regular `close` system call too, which does a complete shutdown, the same as `shutdown(s, SHUT_RDWR)`.

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Java provides shortcuts that combine calls

Example

```
Socket s = new Socket("www.rutgers.edu", 2211)
```

```
C
int s = socket(AF_INET, SOCK_STREAM, 0);

struct sockaddr_in myaddr; /* Initialize address structure */
myaddr.sin_family = AF_INET;
myaddr.sin_addr.s_addr = htonl(NADDR_ANY);
myaddr.sin_port = htons(0);

bind(s, (struct sockaddr *)&myaddr, sizeof(myaddr));

/* Look up the server's address
struct hostent *hp; /* host information */
struct sockaddr_in servaddr; /* server address */
memset((char *)&servaddr, 0, sizeof(servaddr));
servaddr.sin_family = AF_INET;
servaddr.sin_port = htons(2211);
hp = gethostbyname("www.rutgers.edu");

if (connect(fd, (struct sockaddr *)&servaddr, sizeof(servaddr)) < 0) {
    /* connect failed */
}
```

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Using sockets in Java

- **java.net** package
 - **Socket** class
 - Deals with sockets used for TCP/IP communication
 - **ServerSocket** class
 - Deals with sockets used for accepting connections
 - **DatagramSocket** class
 - Deals with datagram packets (UDP/IP)
- Both **Socket** and **ServerSocket** rely on the **SocketImpl** class to actually implement sockets
 - But you don't have to think about that as a programmer

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Create a socket for listening: server

Server:

- `create`, `name`, and `listen` are combined into one method
- **ServerSocket** constructor

```
ServerSocket svc = new ServerSocket(80, 5);
```

port backlog

Several other flavors (see API reference)

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1. Server: create a socket for listening

```
Server Socket svc = new ServerSocket(80, 5);
```

Send HTTP request message to get a page

Receive HTTP request message
Process HTTP request

Send HTTP response message

Receive HTTP response message
Display a page

time

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Server: wait for (accept) a connection

- **accept** method of **ServerSocket**
 - block until connection arrives
 - return a **Socket**

```
ServerSocket svc = new ServerSocket(80, 5);
Socket req = svc.accept();
```

This is a new socket for this "connection"

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2. Server: wait for a connection (blocking)

```

Server Socket svc = new ServerSocket(80);
Socket req = svc.accept();
    
```

Block until an incoming connection comes in

Client: web browser
Server: web server

Send HTTP request message to get a page → Receive HTTP request message
Process HTTP request
Send HTTP response message → Receive HTTP response message
Display a page

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Create a socket: client

Client:

- create, name, and connect operations are combined into one method
- **Socket** constructor

```

Socket s = new Socket(host, port);
    
```

Several other flavors (see API reference)

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3. Client: connect to server socket (blocking)

```

Socket s = new Socket("pk.org", 80);
    
```

Blocks until connection is set up

Client: web browser
Server: web server

Send HTTP request message to get a page → Receive HTTP request message
Process HTTP request
Send HTTP response message → Receive HTTP response message
Display a page

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3a. Connection accepted

```

Server Socket svc = new ServerSocket(80, 5);
Socket req = svc.accept();
    
```

Client: web browser
Server: web server

Socket s = new Socket("pk.org", 80);
Connection is established → Connection is accepted

Send HTTP request message to get a page → Receive HTTP request message
Process HTTP request
Send HTTP response message → Receive HTTP response message
Display a page

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Exchange data

- Obtain InputStream and OutputStream from Socket
 - layer whatever you need on top of them
 - e.g. DataInputStream, PrintStream, BufferedReader, ...

Example:

client

```

DataInputStream in = new DataInputStream(s.getInputStream());
PrintStream out = new PrintStream(s.getOutputStream());
    
```

server

```

DataInputStream in = new BufferedReader(
    new InputStreamReader(req.getInputStream()));
String line = in.readLine();
DataOutputStream out = new DataOutputStream(
    req.getOutputStream());
out.writeBytes(mystring + '\n')
    
```

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4. Perform I/O (read, write)

```

Server Socket svc = new ServerSocket(80, 5);
Socket req = svc.accept();
InputStream s_in = s.getInputStream();
OutputStream s_out = s.getOutputStream();
InputStream r_in = req.getInputStream();
OutputStream r_out = req.getOutputStream();
    
```

Client: web browser
Server: web server

Send HTTP request message to get a page → Receive HTTP request message
Process HTTP request
Send HTTP response message → Receive HTTP response message
Display a page

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Close the sockets

Close input and output streams first, then the socket

```

client:
try {
    out.close();
    in.close();
    s.close();
} catch (IOException e) {}

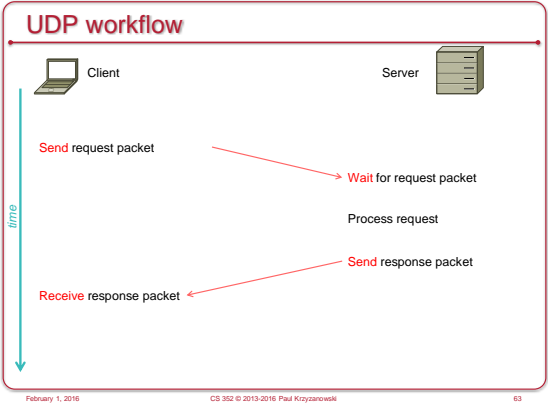
server:
try {
    out.close();
    in.close();
    req.close(); // close connection socket
    svc.close(); // close ServerSocket
} catch (IOException e) {}
    
```

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TCP vs. UDP sockets

- TCP ("stream sockets")
 - Requires a connection (**connection-oriented**)
 - Dedicated socket for accepting connections
 - Communication socket provides a bi-directional link
 - Byte-stream: no message boundaries
- UDP ("datagram sockets")
 - **Connectionless**: you can just send a message
 - Data send in discrete packets (messages)

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Send a packet

```

/* read a line from the user */
BufferedReader user_input = new BufferedReader(new InputStreamReader(System.in));
String line = user_input.readLine();

/* convert it to an array of bytes */
byte[] out_data = line.getBytes();

/* create a datagram socket */
DatagramSocket s = new DatagramSocket();

InetAddress addr = InetAddress.getByName("test.pk.org"); /* look up IP address */
int port = 1234; /* port number */

/* construct the packet */
DatagramPacket out_packet = new DatagramPacket(data, data.length, addr, port);

/* send it out on the socket */
s.send(out_packet);
    
```

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Receive a packet

```

byte in_buf[] new byte[1500];
int port = 4321; /* port number on which we want to receive data */

/* create a datagram socket */
DatagramSocket s = new DatagramSocket(port);

/* create the packet for receiving the data */
DatagramPacket in_packet = new DatagramPacket(in_buf, in_buf.length);

/* get the packet from the socket */
s.receive(in_packet);

System.out.println(
    "received data [" + new String(in_packet.getData(), 0, in_packet.getLength()) + "]" +
    " from address: " + in_packet.getAddress() +
    " port: " + in_packet.getPort());
    
```

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Concurrency & Threads

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Threads

- Designed to support multiple flows of execution **in one process**
- Each thread is scheduled by the operating system's scheduler
- Each thread has its own stack
 - Local variables are local to each thread
- Shared heap
 - Global and static variables and allocated memory are shared
- Multi-core processors make threading attractive
 - Two or more threads can run at the same time

Appeal of threads

- One process can handle multiple requests at the same time
 - Some threads may be blocked
 - Does not affect the threads that have work to do
- User interactivity possible even if certain events block
 - Examples:
 - disk reads
 - wait for network messages
 - count words
 - justify text
 - check spelling

Java Threads

- Create a class that **extends Thread** or **implements Runnable**
- Instantiate this class *or* a Thread to run this Runnable
- When the **run** method is invoked, it starts a new thread of execution
 - After the caller returns, the run method is still running ... as a separate thread
 - Call **join** to wait for the run method to terminate (return)

Java Threads example

```

/* Worker defines the threads that we'll create */
Class Worker extends Thread {
    Worker(...) { // constructor
    }
    public void run() {
        /* thread's work goes here */
        /* thread exits when run() is done */
    }
}

/* other code to start thread */
Worker T = new Worker(); // constructor

T.start(); // start new thread in run method
          // original thread keeps running ...

T.join(); // wait for T's thread to finish.
    
```

Java Threads

Example of threads in a server

- Main thread
 - Waits for new requests from clients
 - After an **accept**, create a **worker thread** to handle the socket connection for that client
- Worker thread handles the request for the client
 - Returns when done – thread disappears

Example of threads in a server

This example shows threads with "implements Runnable"

```

for(;;){
    Socket r = ss.accept(...); /* wait for a new connection */
    doWork worker = new doWork(r); /* create the object */
    Thread t = new Thread(worker); /* create the thread */
    t.start(); /* start running it */
    /* ... and loop back to wait for the next connection */
}

public class doWork implements Runnable {
    private Socket sock;

    doWork(Socket sock) {
        this.sock = sock;
    }

    public void run() { /* here's where the work is done */
        DataInputStream in = new DataInputStream(sock.getInputStream());
        PrintStream out = new PrintStream(server.getOutputStream());
        /* do the work */
        sock.close();
    }
}
    
```

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Threads allow concurrent access

- Threads allow shared access to shared data
- If two threads access the the same data at the same time, results can be undefined

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Race conditions

A **race condition** is a bug:

- The outcome of concurrent threads is unexpectedly dependent on a specific sequence of events

Example

- Your current bank balance is \$1,000
- Withdraw \$500 from an ATM machine while a \$5,000 direct deposit is coming in

<u>Withdrawal</u>	<u>Deposit</u>
• Read account balance	• Read account balance
• Subtract \$500	• Add \$5,000
• Write account balance	• Write account balance

Possible outcomes:
 Total balance = \$5500 (✓), \$500 (X), \$6000 (X)

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Synchronization

- Synchronization: techniques to avoid race conditions
 - Prevent concurrent access
- Operating systems may give us:
 - Semaphores, messages, condition variables, event counters
- Synchronization in Java
 - Add the keyword **synchronized** to a method
 - JVM ensures that at most one thread can execute that method at a time

```

Account {
    double total;
    public synchronized void withdraw(double amount) {
        this.total -= amount;
    }
    public synchronized void deposit(double amount) {
        this.total += amount;
    }
}
    
```

These two methods will never execute concurrently if they're in the same object

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Finer-grain synchronization: blocks

- The **synchronized** keyword provides method-level mutual exclusion
 - Among all methods that are synchronized, only 1 can execute at a time
- Synchronized block: create a mutex for a region

```

Account {
    double total;
    public void withdraw(double amount) {
        synchronized(this.total) {
            this.total -= amount;
        }
    }
    public void deposit(double amount) {
        synchronized(this.total) {
            this.total += amount;
        }
    }
}
    
```

These two blocks will never execute concurrently

this.total becomes a monitor object.

Only one thread can execute in a block synchronized on the same monitor object

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The end

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