

Some Terms

- Ad hoc mode (peer-to-peer mode)
 - No back-end infrastructure is present
 - Hosts have to figure out address assignment, name resolution, and routing among themselves
 - Often no base stations: connectivity directly to hosts and routing via forwarding through hosts



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802.11 LANs

- 802.11 = Wi-Fi
 - Set of standards for wireless local area networking

Standard	Frequency (GHz)	Data rate (max)
802.11	2.4	1-2 Mbps (obsolete)
802.11b	2.4	11 Mbps
802.11a	5	54 Mbps
802.11g	2.4	54 Mbps
802.11n	2.4, 5	72.2 Mbps
802.11ac	5	1.3 Gbps
802.11ad	60	6.9 Gbps (in-room)

5 GHz = 5.1-5.8 GHz 2.4 GHz = 2.5-2.485 GHz

And more... 802.11af, 802.11ah, 802.11aj, 802.11ay

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802.11 LANs

- Base station = access point (AP)
- Basic Service Set (BSS)
 - One or more wireless stations (devices)
 - and one central access point (AP)
- BSSID = MAC address of the AP
- Devices using an AP operate in infrastructure mode
 - AP interconnects with the wired Ethernet infrastructure
- 802.11 devices can also operate in ad hoc mode
 - Communicate with each other directly

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Access Point Identification

- · An access point is assigned
 - A Service Set Identifier (SSID) = textual name for the BSSID
 - A channel number
 - Frequency band is divided into multiple overlapping channels
 - 802.11g/n has 3 non-overlapping channels in the U.S. (1, 6, 11)

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Access Point Discovery & Association

A wireless host (station) needs to associate with one AP

Passive Scanning

- AP periodically sends beacon frames, each containing the AP's SSID & MAC address
- Wireless station scans all channels, searching for beacon frames from any APs

Active Scanning

 Wireless station may also broadcast a probe frame to all APs – iterating through the channels

Selection

- Wireless station selects one access point (often chosen by the user)
- Sends association request frame; receives an association response from AP
- Then send a DHCP discovery message ...

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802.11 MAC Protocol

- · Key differences between Ethernet and 802.11
 - Higher bit-error rates in wireless
 - Ethernet can listen while transmitting; 802.11 cannot
 - · Received signal is weaker than transmitted signal
 - · Receiving station may be receiving signals that the transmitter cannot detect
 - Because Ethernet could listen, it could stop transmission if collision
- What does 802.11 do?
 - Uses Link-layer acknowledgements (ARQ; ack & retransmission)
 - Use CSMA/CA
 - CSMA/CA: Carrier Sense Multiple Access with Collision Avoidance
 - · Random access protocol
 - · Avoid collisions when possible
 - If two stations sense a busy channel, they both enter random backoff

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802.11 MAC Protocol: CSMA/CA

Key idea

- Prevent collisions when they are most likely to occur: when when nodes sense that the channel is clear
- Force nodes to wait a random time, sense, and transmit
- If the channel is busy, the node freezes its timer until it is free
- This reduces the chance that two clients will transmit simultaneously

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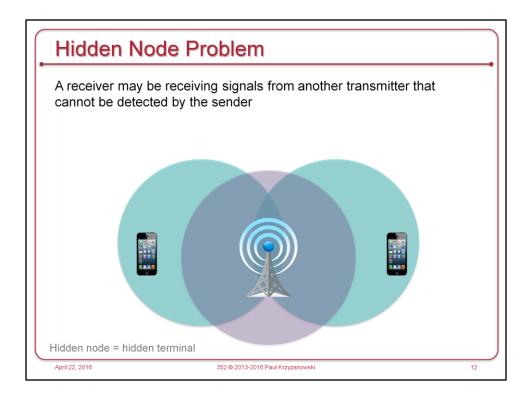
802.11 MAC Protocol: CSMA/CA

- 1. If the channel is idle
 - Wait a short time (Distributed Inter-frame Space, DIFS)
 - Transmit complete frame
- → 2. Else pick a random backoff value using binary exponential backoff
 - Count down this amount when the channel is sensed idle
 - If the channel is busy, the counter does not change
 - 3. When the counter reaches zero (channel must be idle)
 - Transmit the complete frame
 - 4. Wait for an acknowledgement
 - If a receiver receives the frame & CRC is OK,
 - Waits briefly (Short Inter-frame Spacing, SIFS)
 - · Sends back an acknowledgement frame
 - If the transmitter has another frame to send, go to step 2 with new frame
 - If the ACK was not received, increase the backoff value; go to step 2

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802.11 MAC: RTS/CTS

- · Carrier sensing suffers from the hidden node problem
- RTS/CTS: Additional mechanism for sensing in 802.11 (optional)
 - Before sending a frame, send a Request to Send (RTS) frame to AP
 - · Reserves access to the channel
 - · RTS indicates the size of the data frame that will be sent
 - AP responds with a broadcast Clear to Send (CTS) frame
 - · Gives permission to send the frame
 - · Informs other stations not to send anything during that time
 - RTS & CTS frames age generally much shorter than data frames
 - · Minimizes collision
 - RTS/CTS has an overhead
 - Used only for large frames > threshold

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802.11 Frame

- · Similarities to Ethernet frame
 - Same 6-byte MAC addresses
 - Payload
 2312 bytes vs. Ethernet's 1500 bytes, but normally kept ≤ 1500 bytes
 - 32-bit CRC checksum
- · Key difference
 - Ethernet has two address fields: source address & destination address
 - 802.11 has four address fields!
 - · Three addresses are always used
 - · Four are only used for Ad hoc mode
- Also: 802.11n and 802.11ac support optional use of ECC (Low-Density Parity Check codes, LDPC)

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802.11 MAC Addresses

An AP needs to interconnect between the BSS and a wired LAN

Address 1: (wireless destination)

- MAC address of the wireless station that will receive the frame
- If a wireless station transmits, this is the address of the AP
- If an AP is sending to a wireless station, this is the address of the station

Address 2: (wireless source)

- MAC address of the wireless station that transmits the frame
- If a wireless station transmits, this is the address of the station
- If the AP is sending, this is the MAC address of the AP

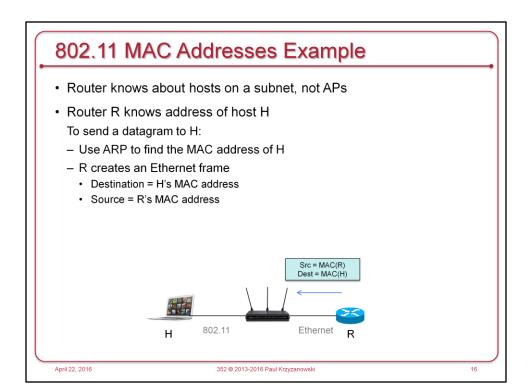
Address 3 (wired destination/source)

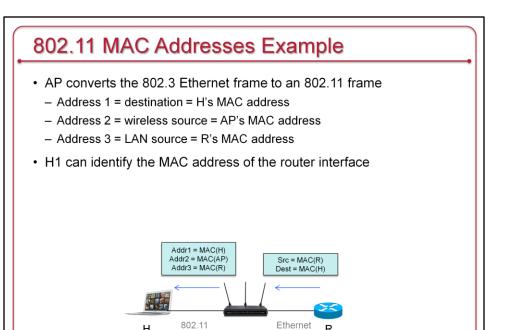
- MAC address of the device on the wired network

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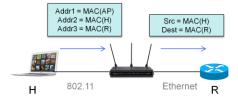


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802.11 MAC Addresses Example

- Return datagram from H to R
- H creates an 802.11 frame
 - Address1 = wireless destination = AP's MAC address
 - Address 2 = source = H's MAC address
 - Address 3 = ultimate LAN destination = R's MAC address
- The AP then creates an Ethernet MAC frame for
 - Source address = H's MAC address
 - Destination address = R's MAC address



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ARQ Protocol & Retransmissions

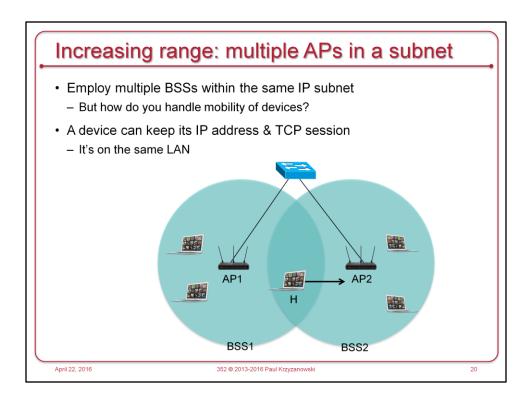
ARQ = Automatic Repeat Request

- Unlike Ethernet, 802.11 uses an ARQ protocol
 - We saw that ACKs can get lost, resulting in retransmissions
 - Retransmissions → duplicate packets
- 802.11 has a sequence number in its MAC header
 - Allows a receiver to distinguish duplicate packets from new packets

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Increasing range: multiple APs in a subnet

- Host migration
 - A host detects a weakening signal from its associated AP (AP1)
 - Scans for an AP with a stronger signal
 - Detects an AP with the same SSID but a stronger signal (AP2)
 - Dissociates with AP1 and associates with AP2

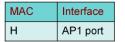
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Increasing range: multiple APs in a subnet

- · What about the switch?
 - Switches are self-learning
 - Switch has an entry in its forwarding table
 - · Associates H'a MAC address with the switch interface to AP1
 - When H associates with BSS2:
 - AP2 will send a broadcast Ethernet frame with H's source address to the switch
 - · The switch will update its forwarding table



MAC Interface
H AP2 port

initial forwarding table

after forged broadcast from AP2

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802.11 Power Management

- · A transceiver on a node can switch between sleep and wake modes
- A node tells its AP that it will go to sleep
 - Sets a power management bit in the 802.11 MAC header
 - Timer in the transceiver is set to wake before the AP is scheduled to send its beacon frame (typically every 100 ms)
- · Frame buffering
 - AP knows that a node went to sleep
 - · Any frames for the node are stored at the AP
 - · Beacon frame contains a list of nodes with buffered frames
 - If no frames to receive, the node goes back to sleep
 - · Otherwise, it requests the buffered frames by sending a polling message
- This can achieve 99%+ sleep times

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Bluetooth

- Bluetooth = IEEE 8002.15.1 → designed as cable replacement
- Short-range, low-power, relatively low-speed (up to 4 Mbps), cheap
- Media
 - 2.4 GHz band 625 μs time slots TDM network access
 - Sender transmits on one of 79 channels
 - Frequency Hopping Spread Spectrum (FHSS)
- Ad hoc network
 - No access point
 - Up to 8 active devices (255 "parked" devices)
 - One designated as a master others are slaves
 - · Master can transmit in each odd-numbered slot
 - · Slaves transmit only after master granter permission and only to the master

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Wide Area Mobility: Cellular Networking

- Home Network
 - Permanent device address
- Foreign Network
 - Foreign agent responsible for
 - Care-of-Address (COA) = foreign address
 - · Can be obtained via DHCP on the foreign network
 - Informing Home Agent of the node's current foreign address

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Mobile Routing: Indirect Routing

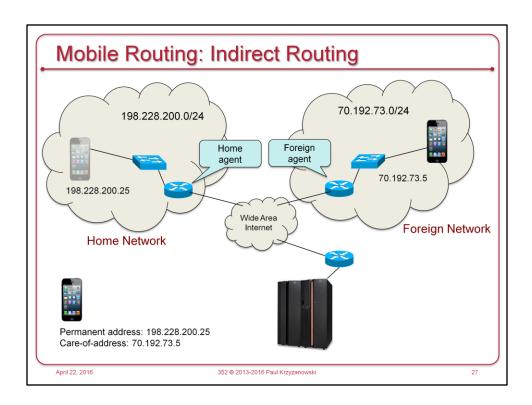
- · To the mobile node
 - Address datagrams to mobile node's permanent address
 - Datagrams get routed to the home network
 - Home agent
 - · Tracks COAs
 - Intercepts datagrams for nodes residing on foreign networks
 - Encapsulates datagrams & forwards them to the foreign agent
 - Outer datagram is addressed to the foreign agent
 - Inside datagram is the original datagram
 - Foreign agent extracts the encapsulated datagram & forwards to node
- From the mobile node
 - Mobile node can send datagrams directly from its permanent address

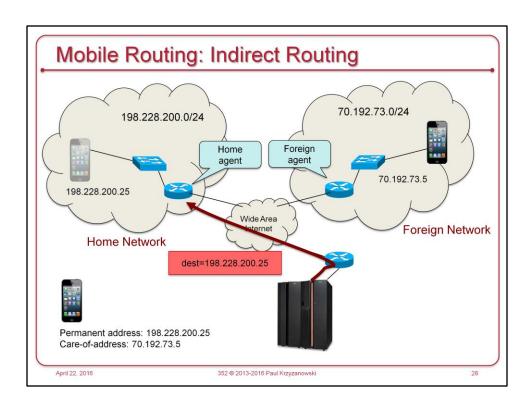
Mobile IP: RFC 5944

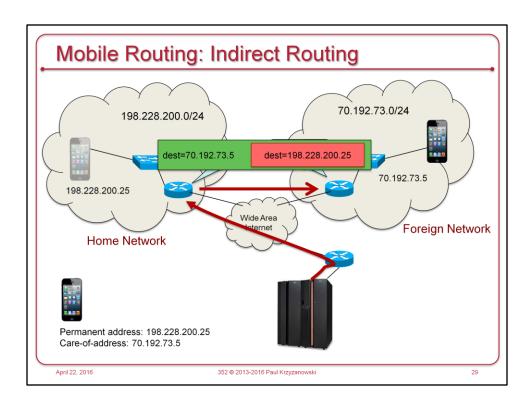
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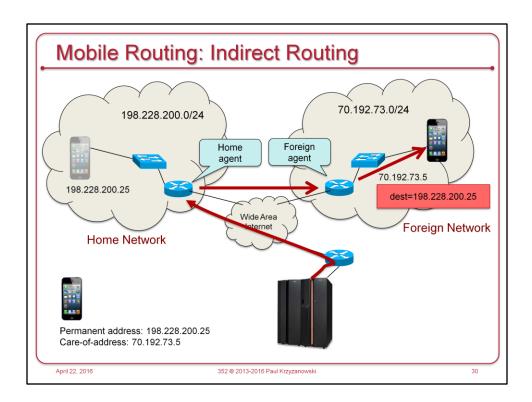
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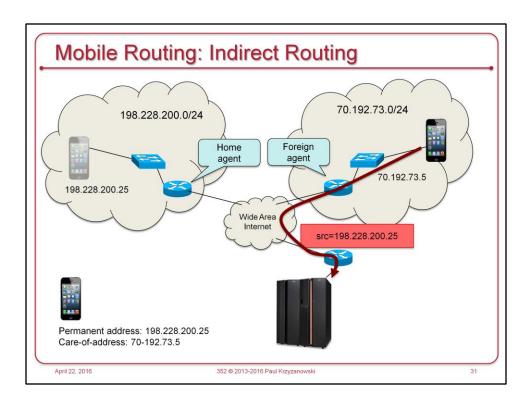
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Mobile Routing: Direct Routing

- Indirect routing suffers from the triangle routing problem
 - Datagrams to the mobile node must be routed through the home node
- Direct Routing
 - Add a Corresponding Agent to the sender's network
 - Learns the care-of-address (COA) of the mobile node
 - · Query home agent to find the COA & foreign agent
 - Original foreign agent = anchor foreign agent
 - If the mobile node moves to another foreign network
 - · Mobile node registers with the new foreign agent
 - · New foreign agent tells the anchor foreign agent the new COA
 - Anchor foreign agent encapsulates incoming datagrams and routes them to the new foreign agent (indirect routing)

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